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## Updating the regularization parameter in the adaptive cubic regularization algorithm

Nicholas I. M. Gould<sup>1,2</sup>, Margherita Porcelli<sup>3</sup> and Philippe L. Toint<sup>3</sup>

#### ABSTRACT

The adaptive cubic regularization method (Cartis, Gould & Toint, Math. Programming, DOI: 10.1007/s10107-009-0286-5 & 10.1007/s10107-009-0337-y) has been recently proposed for solving unconstrained minimization problems. At each iteration of this method, the objective function is replaced by a cubic approximation which comprises an adaptive regularization parameter whose role is related to the local Lipschitz constant of the objective's Hessian. We present new updating strategies for this parameter based on interpolation techniques, which improve the overall numerical performance of the algorithm. Numerical experiments on large nonlinear least-squares problems are provided.

Keywords: unconstrained optimization, cubic regularization, numerical performance. AMS classification: 49J52, 49M37, 65F22, 65K05, 90C26, 90C30, 90C55.

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## 1 Introduction

We consider the unconstrained minimization problem

$$\min_{x \in \mathbb{R}^n} f(x), \tag{1.1}$$

where f is a twice continuously differentiable function of the variables  $x \in \mathbb{R}^n$ . A simplistic method for solving this problem is to compute an improving step  $s_k$  by minimizing a quadratic Taylor-series model of the objective function around the current iterate  $x_k$ . Unfortunately, it is well-known that an iteration based on this simple idea may not always be well-defined (when the Taylor model is nonconvex), nor converge globally. These drawbacks may be overcome by restricting the model minimization to a trust region containing  $x_k$  [8]. Clearly, trust-region strategies may be considered as regularization techniques because they control the difference between two consecutive iterates by explicitly imposing a restriction on the stepsize.

The main motivation for this paper is a series of recent papers where alternative regularization strategies are introduced [2, 3, 7, 17, 20, 24]. These procedures are based on the minimization of quadratic or cubic models for the objective function in a neighbourhood implicitly defined by a regularization term that penalizes the step length. In particular, the adaptive cubic regularization (ARC) algorithm is proposed in [3] for solving problem (1.1). At each iteration, the objective function is locally replaced by a cubic approximation, in which third- and higher-order Taylor-series terms are replaced by a cubic regularization term, and an adaptive estimation of the local Lipschitz constant of the objective function's Hessian is employed. The method has been shown to have excellent global and local convergence properties and numerical experiments indicate that the new procedure may be competitive with the trust region approach when solving small-scale problems [3]. Additionally, and of theoretical interest, ARC possesses a better worst-case evaluation-complexity bound than its trust-region competitor [5].

The purpose of this paper is twofold. Firstly, we propose alternative updating rules for the regularization parameter of the ARC algorithm which are based on interpolation techniques. In particular, in the trust-region case, the restriction on the stepsize is explicitly imposed by the trust-region constraint. By contrast, in the cubic regularization case the control on the stepsize is nonlinear and is defined implicitly. This suggests a need to design an efficient updating rule for the regularization parameter that is able to control the stepsize in a flexible way.

Secondly, we shall apply these ideas and report on extensive numerical experiments on the solution of large nonlinear least-squares problems, that is problems of the form

$$\min_{x \in \mathbb{R}^n} f(x) = \frac{1}{2} \|h(x)\|_2^2, \tag{1.2}$$

where  $h : \mathbb{R}^n \to \mathbb{R}^m$  is a given continuously differentiable mapping. By limiting our discussion to this problem, we may specialize the models employed in both the ARC and trust-region algorithms to those that are suited to solving nonlinear least-squares

problems, specifically using regularized Gauss-Newton-based models, and consequently to take advantage of the ideas and implementations details proposed in [7] for the solution of large regularized linear least-squares problems. Since we are primarily interested in large problems for which matrix factorization often has prohibitive computational cost, we shall focus on iterative algorithms for the subproblems, particularly on those implemented as part of version 2.4 of the GALAHAD optimization library [16]. Such procedures are based on the minimization of the local model of the objective function over a sequence of (nested) subspaces associated with the Lanczos procedure. As a result, they are especially suited to the large-scale setting and allow us to test the methods on large problems from the CUTEr test collection [15]. In particular, the new updating rules for the regularization parameter of the ARC algorithm are experimentally validated and a comparison with the trust-region algorithm is performed on problem (1.2).

The paper is organized as follows. In Section 2 we review the standard trust-region algorithm and the ARC algorithm for the solution of problem (1.1). New updating rules for the regularization parameter in the ARC algorithm are introduced in Section 3. Section 4 is dedicated to numerical experiments and, finally, in Section 5 we draw some conclusions.

Throughout the paper we use the following notation. The Euclidean  $(\ell_2)$  norm is denoted by  $\|\cdot\|$ , and I represents the identity matrix. Given a sequence of vectors  $\{x_k\}$ , for any generic function h we let  $h_k = h(x_k)$ . Let  $g(x) = \nabla f(x)$  where f is the objective function in (1.1) and let J(x) denote the Jacobian matrix of the residual function h(x) in (1.2). Finally,  $\epsilon_m \approx 10^{-16}$  denotes the relative machine (double) precision.

## 2 The algorithms

In this section, we describe the kth iteration of two globally convergent algorithms for the solution of problem (1.1): the standard trust-region algorithm (e.g. [8]) and the ARC algorithm ([3]).

In the trust-region framework, a quadratic model of f(x) around  $x_k$  is constructed by defining the model of the objective function to be

$$q_k(s) = f_k + g_k^T s + \frac{1}{2} s^T H_k s, \qquad (2.1)$$

where  $H_k$  is a symmetric approximation to the local Hessian  $\nabla_{xx} f_k$ . Then, a trial step  $s_k$  is computed by solving (possibly only approximately) the subproblem

$$\min_{s \in \mathbb{R}^n} \{ q_k(s) : \|s\| \le \Delta_k \}, \tag{2.2}$$

where  $\Delta_k > 0$  is the so-called trust-region radius.

By contrast, assuming that the objective's Hessian  $\nabla_{xx} f$  is globally Lipschitz continuous on  $\mathbb{R}^n$  with Lipschitz constant L, the cubic model used in the ARC algorithm is based

on the bound

$$f(x_{k}+s) = f_{k} + s^{T}g_{k} + \frac{1}{2}s^{T}\nabla_{xx}f_{k}s + \int_{0}^{1}(1-\tau)s^{T}[\nabla_{xx}f(x_{k}+\tau s) - \nabla_{xx}f_{k}]s d\tau$$
  

$$\leq f_{k} + s^{T}g_{k} + \frac{1}{2}s^{T}\nabla_{xx}f_{k}s + \frac{1}{6}L||s||^{3} \stackrel{\text{def}}{=} l_{k}(s), \qquad (2.3)$$

which holds for for all  $s \in \mathbb{R}^n$ . Thus, so long as  $l_k(s_k) < l_k(0) = f_k$ , the new iterate  $x_{k+1} = x_k + s_k$  improves f(x). In [3], a dynamic positive parameter  $\sigma_k$  replaces the Lipschitz constant L/2 and a symmetric approximation  $H_k$  to the local Hessian  $\nabla_{xx} f_k$  is allowed. At each iteration, the cubic model

$$c_k(s) = f_k + s^T g_k + \frac{1}{2} s^T H_k s + \frac{1}{3} \sigma_k ||s||^3, \qquad (2.4)$$

is employed as an approximation to the objective f and the subproblem

$$\min_{s \in \mathbb{R}^n} c_k(s) \tag{2.5}$$

is solved. The parameter  $\sigma_k$  plays a crucial role in the description of the ARC algorithm as it measures the discrepancy between the objective function and its second order Taylor expansion and of the difference between the exact and the approximate Hessian [3].

It is important to note that the restriction on stepsize is explicitly imposed by the trust-region constraint in the trust-region case, while stepsize control is defined implicitly, indeed nonlinearly, in the cubic case. In fact, a step  $s_k$  derived by reducing (2.5) is always bounded [3, Lem.2.2] by

$$||s_k|| \le 3 \max\left[\frac{||H_k||}{\sigma_k}, \sqrt{\frac{||g_k||}{\sigma_k}}\right].$$

Such a bound suggests that the regularization parameter  $\sigma_k$  for the ARC algorithm may loosely be interpreted as the reciprocal of the trust-region radius  $\Delta_k$ . This observation in turn suggests choosing updating rule for the parameter  $\sigma_k$  by analogy with the trust-region case. In a standard trust-region scheme, the trust-region radius may be enlarged if there is a sufficient decrease in f(x), computed by some measure of the relative objective changes, and it is reduced otherwise. In the regularization case, the parameter  $\sigma_k$  is decreased if there is a sufficient agreement between the objective function and the model, but increased or left unchanged otherwise.

In both algorithms, the agreement between the model and the objective function is given by the standard ratio of the achieved to the predicted reduction, and the size of this ratio is used to decide whether or not to accept the trial step and to change the regularization parameter. This ratio takes the form

$$\rho_q(s_k) = \frac{f_k - f(x_k + s_k)}{q_k(0) - q_k(s_k)},$$
(2.6)

in the trust-region case, and

$$\rho_c(s_k) = \frac{f_k - f(x_k + s_k)}{c_k(0) - c_k(s_k)},\tag{2.7}$$

in the the cubic regularization case, where the models  $q_k$  and  $c_k$  are defined in (2.1) and (2.4) respectively. Without ambiguity, let  $\rho(s)$  represent both  $\rho_c(s)$  and  $\rho_q(s)$ , and let  $\eta_1$ ,  $\eta_2$  be constants such that  $0 < \eta_1 < \eta_2 < 1$ . We say that the iteration k is very successful if  $\rho(s_k) \ge \eta_2$ , successful if  $\rho(s_k) \in [\eta_1, \eta_2)$ , unsuccessful otherwise. When it is useful to distinguish the case  $\rho(s_k) < 0$  within the unsuccessful case, we refer to a very unsuccessful iteration.

The general framework of the methods described so far is presented in Algorithm 2.1. The string METHOD denotes the name of the method, i.e. it is either 'TRUST-REGION' or 'ARC'. Sections 2.1 and 3 give further insight into Steps 1 and 4.

#### 2.1 Computing a trial step

Step 1 of Algorithm 2.1 leaves substantial implementation freedom, which may be used according to context. The focus of this paper is on the case where matrix factorizations of the Hessian matrix are not feasible, implying that iterative methods for computing a trial step are needed. We consider the class of subspace minimization methods, i.e. methods that find an approximate solution by solving a sequence of minimization problems with the additional constraint that s is contained in a subspace. This class may be divided into two subclasses depending on the construction of the sequence of subspaces. The first consists of expanding subspaces methods. The Conjugate Gradient (CG) method belongs to this subclass as it may be viewed as a subspace minimization method for finding an unconstrained minimizer of a strictly convex quadratic function, where, at each successive iteration, the quadratic function is minimized by restricting the variable to a sequence of nested Krylov subspaces. In [3, 14], methods based on this approach have been proposed for solving the regularized cubic problem (2.5) and the trust-region problem (2.2), respectively. The second subclass comprises low-dimensional subspace methods, i.e. methods that always generate subspaces of low-dimension. Such methods have been proposed in literature only for solving problem (2.2) and differ in the choice of the subspaces [11, 12, 18, 19]. In order to apply the same subspace approach to both the trust-region and the cubic case, we consider the former subclass of methods to perform Step 1.

Consider the nonlinear least-squares problem (1.2). At the current iterate  $x_k$ , the exact Hessian of the objective function f has the form

$$\nabla_{xx} f_k = J_k^T J_k + S_k,$$

where  $S_k$  contains the second-order information on the residual. If  $S_k$  is small, it is reasonable to consider the first order approximation  $H_k = J_k^T J_k$ . This is the case, for instance, in

#### Algorithm 2.1: Generic trust-region/cubic regularization method

An initial point  $x_0$  as well as constants  $0 < \eta_1 < \eta_2 < 1$  and  $\gamma > 1$  are given. If METHOD = 'TRUST-REGION', set the initial radius  $\Delta_0 > 0$  and the constants  $\tau_1, \tau_2$  such that  $0 \le \tau_1 \le \tau_2 \le 1$ . Else set the initial regularization parameter  $\sigma_0 > 0$  and the constants  $\nu_1, \nu_2$  such that  $1 < \nu_1 \le \nu_2$ .

For  $k = 0, 1, \ldots$ , until convergence,

- Step 1: Trial step computation. If METHOD = 'TRUST-REGION', compute  $s_k$  as an (approximate) solution of problem (2.2). Else, compute  $s_k$  as an (approximate) solution of problem (2.5).
- Step 2: Step acceptance. If METHOD = 'TRUST-REGION', compute  $\rho(s_k) = \rho_q(s_k)$  as in (2.6). Else, compute  $\rho(s_k) = \rho_c(s_k)$  as in (2.7).

If  $\rho(s_k) \ge \eta_1$ , let  $x_{k+1} = x_k + s_k$ ; otherwise let  $x_{k+1} = x_k$ .

Step 4: Regularization parameter update. If METHOD = 'TRUST-REGION' set

$$\Delta_{k+1} \in \begin{cases} [\Delta_k, \infty) & \text{if } \rho(s_k) \ge \eta_2 \quad \text{[very successful iteration]} \\ [\tau_2 \Delta_k, \Delta_k] & \text{if } \rho(s_k) \in [\eta_1, \eta_2) \quad \text{[successful iteration]} \\ [\tau_1 \Delta_k, \tau_2 \Delta_k] & \text{otherwise} \quad \text{[unsuccessful iteration]} \end{cases} . (2.8)$$

Else set

$$\sigma_{k+1} \in \begin{cases} (0, \sigma_k] & \text{if } \rho(s_k) \ge \eta_2 & \text{[very successful iteration]} \\ [\sigma_k, \nu_1 \sigma_k] & \text{if } \rho(s_k) \in [\eta_1, \eta_2) & \text{[successful iteration]} \\ [\nu_1 \sigma_k, \nu_2 \sigma_k] & \text{otherwise} & \text{[unsuccessful iteration]} \end{cases}$$
(2.9)

a neighborhood of a zero residual solution of problem (1.2), [10]. Using the approximation  $H_k = J_k^T J_k$ , the quadratic model in (2.1) takes the form

$$q_k(s) = \frac{1}{2} \|J_k s + h_k\|^2, \qquad (2.10)$$

which is the Gauss-Newton model for f, and the cubic model in (2.4) becomes

$$c_k(s) = \frac{1}{2} \|J_k s + h_k\|^2 + \frac{\sigma_k}{3} \|s\|^3,$$
(2.11)

yielding a Gauss-Newton model regularized by a cubic term.

Procedures have been proposed in [7] to solve the subproblems (2.2) and (2.5) in the special case where the models are given in (2.10) and (2.11) respectively. The core component of these procedures is the Golub and Kahan bi-diagonalization process [13] that

generates orthonormal basis of a sequence of expanding subspaces  $\{\mathcal{V}_j\}_{j\geq 1}$ . Let  $V_j \in \mathbb{R}^{n \times j}$ be the orthonormal matrix whose columns span  $\mathcal{V}_j$ . The solutions of problems (2.2) and (2.5) are found by computing the sequence of minimizers  $y_j$  of the reduced problems

$$\min_{y \in \mathbb{R}^j} \{ q_k(V_j y) : \|y\| \le \Delta_k \},$$

$$(2.12)$$

and

$$\min_{y \in \mathbb{R}^j} c_k(V_j y), \tag{2.13}$$

respectively, increasing the dimension j of the subspaces until  $s_j = V_j y_j$  is sufficiently accurate. At that point, the step  $s_k$  in the full space is taken as the last computed  $s_j$  [7].

It is interesting to note, that if the LSQR algorithm [21] is used to solve the unconstrained problem  $\min_{s} q_k(s)$ , a basis of the Krylov subspaces

$$\mathcal{K}_j = \left\{ (J_k^T J_k)^i J_k^T h_k \right\}_{i=0}^{j-1},$$

is given by the columns of  $V_j$ . Due to the equivalence between the LSQR and CG methods, the sequence  $s_i$  generated by LSQR has the favorable property to be monotonically increasing in norm [23]. Thus, either LSQR finds a solution in the interior of the trustregion, or finds an iterate  $s_j$  s.t.  $||s_{j-1}|| \leq \Delta_k < ||s_j||$  and in this case we may conclude that the solution of the problem (2.2) lies on the boundary of the trust-region. When this happens two alternative strategies can be followed: either the so-called Steihaug-Toint point  $[8, \S7.5.1]$  is computed or a solution on the boundary is computed to any prescribed accuracy. The Steihaug-Toint strategy interpolates the last interior iterate  $s_{i-1}$  with the newly discovered exterior one  $s_i$  to find the boundary point between them. The resulting step has the favorable property that the optimal decrease of  $q_k$  at the exact solution of the trust-region problem (2.2), is no more than twice that achieved at the Steihaug-Toint point (see [25] or [8, Thm.7.5.9]). On the negative side however, it makes no attempt to find a constrained solution with prescribed accuracy. A more refined strategy solves a sequence of constrained reduced problems (2.12) increasing j until  $s_j$  is sufficiently accurate [7]. Note that this strategy specializes to problem (2.12) the GLTR method [14] for the general trust-region problem (2.2) in which the CG method is used as long as the iterates are in the interior of the trust-region and the expanding subspaces are defined by the Lanczos vectors.

## **3** Updating rules for the regularization parameters

Because of its central role, the definition of a procedure to update the regularization parameters at Step 4 of Algorithm 2.1 may have a crucial influence on its overall performance. In this section, we first review two established updating strategies for the trust-region radius  $\Delta_k$  and then propose new strategies for the parameter  $\sigma_k$  for the ARC algorithm. Clearly, the rule (2.8) in Algorithm 2.1 leaves considerable flexibility. A simple and reasonable choice is to select

$$\Delta_{k+1} = \begin{cases} \max\{\gamma_2 \| s_k \|, \Delta_k\} & \text{if } \rho_q(s_k) \ge \eta_2 \\ \Delta_k & \text{if } \rho_q(s_k) \in [\eta_1, \eta_2) \\ \gamma_1 \| s_k \| & \text{otherwise} \end{cases} \text{ [unsuccessful iteration], } (3.1)$$

where  $\gamma_1$  and  $\gamma_2$  are constants such that  $0 < \gamma_1 < 1 \leq \gamma_2$ , but further refinements are possible using interpolation techniques in the unsuccessful case. If  $\rho_q(s_k)$  is negative, the agreement between the model and the objective function is extremely poor and some drastic action might be warranted. In this case, we presume for simplicity that  $s_{k+1}$  will be aligned with  $s_k$  and we compute a trust-region radius small enough to ensure that the new step gives at least a successful iteration [8, Chapter 17]. To compute such a radius, we consider a step of the form  $\alpha s_k$  with  $\alpha > 0$  and we set  $\Delta_{k+1} = \alpha_{\eta}^{bad} \Delta_k$  where  $\alpha_{\eta}^{bad}$  solves  $\rho_q(\alpha s_k) = \eta$ , which is equivalent to the scalar nonlinear equation

$$f_k - f(x_k + \alpha s_k) = \eta(q_k(0) - q_k(\alpha s_k)),$$
(3.2)

with  $\eta \in [\eta_1, 1)$  and  $\eta_1$  as given in Algorithm 2.1. To avoid the expense of computing the extra function value  $f(x_k + \alpha s_k)$  and to simplify the solution of (3.2), the scalar function  $\hat{f}(\alpha) = f(x_k + \alpha s_k), \alpha > 0$  is replaced by a quadratic interpolating polynomial for  $\hat{f}$ . The polynomial  $t_f(\alpha)$  such that  $t_f$  and  $t'_f$  agree with  $\hat{f}$  and  $\hat{f}'$  at 0, and  $t_f(1) = \hat{f}(1) = f(x_k + s_k)$ , is given by

$$t_f(\alpha) = f_k + g_k^T s_k \alpha + \left( f(x_k + s_k) - f_k - g_k^T s_k \right) \alpha^2.$$

Substituting this value for  $f(x_k + \alpha s_k)$  into (3.2) and solving for  $\alpha$ , yields the value of  $\alpha_{\eta}^{bad}$  given by

$$\alpha_{\eta}^{bad} = \frac{(1-\eta)g_k^T s_k}{(1-\eta)(f_k + g_k^T s_k) + \eta q_k(s_k) - f(x_k + s_k)}.$$
(3.3)

We may therefore modify (3.1) to use this information and obtain the more sophisticated rule

$$\Delta_{k+1} = \begin{cases} \max\{\gamma_2 \| s_k \|, \Delta_k\} & \text{if } \rho_q(s_k) \ge \eta_2 & \text{[very successful iteration]}, \\ \Delta_k & \text{if } \rho_q(s_k) \in [\eta_1, \eta_2) & \text{[successful iteration]}, \\ \gamma_1 \| s_k \| & \text{if } \rho_q(s_k) \in [0, \eta_1) & \text{[unsuccessful iteration]}, \\ \min\{\gamma_1 \| s_k \|, \max\{\gamma_3, \alpha_\eta^{bad}\} \Delta_k\} & \text{otherwise} & \text{[very unsuccessful iteration]}, \end{cases}$$

$$(3.4)$$

where  $\alpha_{\eta}^{bad}$  is given by (3.3) and the constants  $\gamma_1, \gamma_2, \gamma_3$  are such that  $0 < \gamma_3 < \gamma_1 < 1 \le \gamma_2$ [8].

Let us now consider the ARC framework with this in mind. The updating rule proposed in [3] aims to try to reduce the model rapidly to match the Newton model once convergence sets in, while maintaining some regularization before the asymptotic behaviour. The rule used in the reported experiments was

$$\sigma_{k+1} = \begin{cases} \max\{\min\{\sigma_k, \|g_k\|\}, \epsilon_m\} & \text{if } \rho_c(s_k) \ge \eta_2 \\ \sigma_k & \text{if } \rho_c(s_k) \in [\eta_1, \eta_2) \\ \gamma \sigma_k & \text{otherwise} \end{cases} \text{ [unsuccessful iteration], and } (3.5)$$

with  $\gamma \geq 1$ . Clearly, the relationship between the step length and the regularization parameter in (3.5) is not as simple as in the updating rules (3.1) for the trust-region case and the control of the first by the second is performed implicitly.

To relate the step size and the parameter  $\sigma_k$  in a more direct way, we now present an alternative strategy for updating  $\sigma_k$  in the spirit of the interpolation procedures used with the trust-region scheme. Specifically, we try to ensure, in the very unsuccessful case, that the next iterate gives at least a successful iteration. In the very successful case we may also exploit the overestimation property (2.3) measuring at each iteration the gap between the current objective function value  $f(x_k + s_k)$  and the current model value  $c_k(s_k)$  and reduce  $\sigma_k$  in order to decrease this gap (cf. [17, 24]). In particular, given the current  $x_k, \sigma_k$  and  $s_k$ , we presume, as above, that  $s_{k+1}$  is of the form  $\alpha s_k, \alpha > 0$  and compute the value  $\sigma_{k+1}$  to ensure suitable conditions on  $\alpha s_k$ .

As in the trust-region case, we avoid the need to compute the value of  $f(x_k + \alpha s_k)$ by using instead a suitable interpolating approximation. The interpolating cubic function  $p_f(\alpha)$ ,  $\alpha \ge 0$  we use here is built by requiring that  $p_f(0) = f_k$ ,  $p'_f(0) = g_k^T s_k$ ,  $p''(0) = s_k^T H_k s_k$  and  $p_f(1) = f(x_k + s_k)$ , and hence takes the form

$$p_f(\alpha) = f_k + g_k^T s_k \alpha + \frac{1}{2} s_k^T H_k s_k \alpha^2 + p_{f_3} \alpha^3, \qquad (3.6)$$

where

$$p_{f_3} = f(x_k + s_k) - q_k(s_k).$$
(3.7)

The quadratic model (2.1) along the direction  $s_k$  may be written as

$$q(\alpha) = f_k + g_k^T s_k \alpha + \frac{1}{2} s_k^T H_k s_k \alpha^2, \qquad (3.8)$$

while its regularized cubic counterpart (2.4) is

$$c(\alpha, \sigma) = q(\alpha) + \frac{\sigma \|s_k\|^3}{3} \alpha^3.$$
 (3.9)

We now define the current overestimation gap  $\chi_k^f$  to be

$$\chi_k^f = c_k(s_k) - f(x_k + s_k).$$
(3.10)

Note that the model  $c_k$  at  $s_k$  overestimates  $f(x_k + s_k)$ , i.e.  $\chi_k^f \ge 0$ , if and only if  $\rho_c(s_k) \ge 1$ .

Consider the very successful  $(\chi_k^f > 0)$  case first, in which case the regularization parameter should be decreased. If the current gap  $\chi_k^f$  is large enough, we aim at reducing it by a factor  $\beta \in (0, 1)$ . Assume first that  $f(x_k + s_k) \ge q_k(s_k)$ . Remembering that the next step should minimize the cubic model (in particular along  $s_k$ ), we thus search for  $\alpha$  and  $\sigma$  such that

$$c(\alpha, \sigma) - p_f(\alpha) = \beta \chi_k^f$$
 and (3.11)

$$\frac{a}{d\alpha}c(\alpha,\sigma) = 0, \tag{3.12}$$

 $c(\alpha, \sigma)$  and  $p_f(\alpha)$  given in (3.9) and (3.6). It follows from (3.11) that

$$\sigma = 3 \frac{\beta \chi_k^f + p_{f_3} \alpha^3}{\alpha^3 \|s_k\|^3} \equiv \sigma_k + 3 \frac{\chi_k^f}{\|s_k\|^3} \left(\frac{\beta - \alpha^3}{\alpha^3}\right),$$
(3.13)

and substituting (3.13) into (3.12), we find that the required  $\alpha$  satisfies the cubic scalar equation

$$3\beta\chi_k^f + g_k^T s_k \alpha + s_k^T H_k s_k \alpha^2 + 3p_{f_3} \alpha^3 = 0.$$
(3.14)

Thus, we determine the root  $\alpha$  of (3.14) which exceeds  $\sqrt[3]{\beta}$  by the least (if there is such a root) and recover  $\sigma_{k,\beta}^*$  from (3.13). If there is no such  $\alpha$ , or if  $\alpha$  is too large, we simply reduce  $\sigma_k$  by a factor  $\delta_1 \in (0, 1)$ .

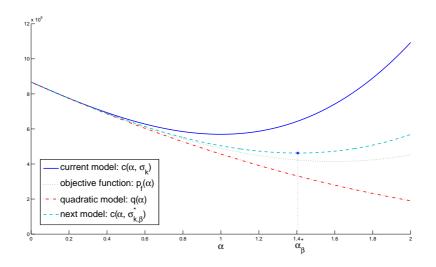


Figure 3.1: Very successful iteration and  $f(x_k + s_k) \ge q_k(s_k)$ .

In Figures 3.1–3.3 the current cubic model  $c(\alpha, \sigma_k)$ , the approximated objective function  $p_f(\alpha)$ , the quadratic model  $q(\alpha)$  and the next cubic model  $c(\alpha, \sigma_{k,\beta}^*)$  are plotted. Figure 3.1 represents an example where the k-th iterate is very successful and  $f(x_k + s_k) \ge q_k(s_k)$ . In this example,  $\beta = 0.5$  and equation (3.14) has two positive roots. The largest one  $(\alpha_{\beta}^*)$  in the figure) is larger than  $\sqrt[3]{\beta} \approx 0.7937$  and gives  $\sigma_{k,\beta}^*$  such that  $\sigma_{k,\beta}^* < \sigma_k$ .

Consider now the case where  $f(x_k + s_k) < q_k(s_k)$ . If we attempt to solve the system (3.11)–(3.12), i.e. try to reduce the quantity  $c_k(s_k) - f(x_k + s_k)$  by a factor  $\beta$ , we might reduce this gap too much, leading to undesirable value of the new  $\sigma$ . Figure 3.2-(a) illustrates the typical situation: in this example ( $\beta = 0.5$ ), equation (3.14) only has one positive solution ( $\alpha^{\chi_k^f} \approx 0.745$  in the figure), but it is smaller than  $\sqrt[3]{\beta}$  so that the corresponding  $\sigma_{k,\beta}^*$  computed by (3.13) is larger than the current  $\sigma_k$ . To avoid this undesirable situation, we instead attempt to reduce the following gap

$$\chi_k^q = c_k(s_k) - q_k(s_k), \tag{3.15}$$

and search for  $\alpha$  and  $\sigma$  such that

$$c(\alpha, \sigma) - q(\alpha) = \beta \chi_k^q$$
 and (3.16)

$$\frac{d}{d\alpha}c(\alpha,\sigma) = 0, \qquad (3.17)$$

with  $c(\alpha, \sigma)$  and  $q(\alpha)$  given in (3.9). Computing  $\sigma$  from (3.16), we then find that

$$\sigma = 3 \frac{\beta \chi_k^q}{\alpha^3 \|s_k\|^3} \equiv \frac{\beta}{\alpha^3} \sigma_k, \qquad (3.18)$$

and substituting (3.18) in (3.17) yields that  $\alpha$  solves the quadratic scalar equation

$$3\beta\chi_k^q + g_k^T s_k \alpha + s_k^T H_k s_k \alpha^2 = 0. \tag{3.19}$$

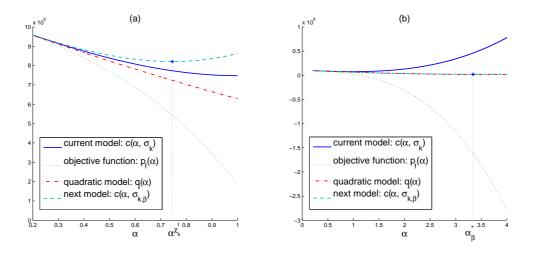


Figure 3.2: Very successful iteration and  $f(x_k + s_k) < q_k(s_k)$ .

As in the previous case, we compute the root of (3.19) which exceeds  $\sqrt[3]{\beta}$  by the least (if such a root exists) and compute the corresponding value  $\sigma_{k,\beta}^*$  using (3.18). Once again, if there is no such  $\alpha$ , or if  $\alpha$  is too large, we simply reduce  $\sigma_k$  by a factor  $\delta_1 \in (0, 1)$ . Figure 3.2-(b) illustrates the same example as in Figure 3.2-(a) but now solving the system (3.16)– (3.17): equation (3.19) has 2 positive roots and one ( $\alpha_{\beta}^*$  in the figure) is larger than  $\sqrt[3]{\beta}$ , so that the corresponding  $\sigma_{k,\beta}^*$  is smaller than  $\sigma_k$ .

Let us now turn to the very unsuccessful case,  $\rho_c(s_k) < 0$ , where we wish to increase the regularization parameter. We proceed as in the trust-region framework simply requiring that  $\alpha s_k$  produces at least a successful iterate. We thus search for  $\alpha$  and  $\sigma$  such that

$$f_k - p_f(\alpha) = \eta(f_k - c(\alpha, \sigma)), \text{ and}$$
 (3.20)

$$\frac{a}{d\alpha}c(\alpha,\sigma) = 0, \tag{3.21}$$

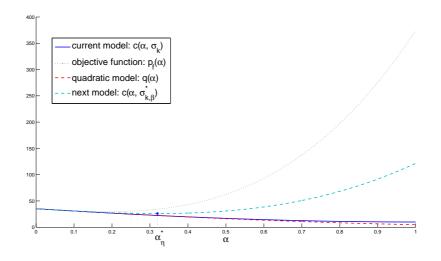


Figure 3.3: Very unsuccessful iteration.

for some  $\eta \in [\eta_1, 1)$ . Computing  $\sigma$  from (3.21) we obtain that

$$\sigma = \frac{-g_k^T s_k - s_k^T H_k s_k \alpha}{\alpha^2 \|s_k\|^3},\tag{3.22}$$

and substituting this expression in (3.20), we find that  $\alpha$  must be a root of the quadratic scalar equation

$$2(3-2\eta)g_k^T s_k + (3-\eta)s_k^T H_k s_k \alpha + 6p_{f_3}\alpha^2 = 0, \qquad (3.23)$$

where  $p_{f_3}$  is positive since  $\rho_c(s_k) < 0$ . The discriminant of the above equation is given by

$$(3-\eta)^2 (s_k^T H_k s_k)^2 - 48(3-2\eta)g_k^T s_k p_{f_3},$$

and as  $\eta < 3/2$ , it is always positive. In this case, the above equation as two roots of opposite sign. If  $\alpha_{\eta}^*$  is the positive one, we then compute  $\sigma_{k,\eta}^*$  from (3.22) with  $\alpha = \alpha_{\eta}^*$ . Figure 3.3 shows an example of this case.

Combining these different cases together, we are now able to state the complete rule for updating the current regularization parameter  $\sigma_k$ : it is described as Algorithm 3.1 on page 14. This algorithm also safeguards against the case where equations (3.14) and (3.19) do not admit a solution larger than  $\sqrt[3]{\beta}$ , or where such a solution exists but may be very much larger than this value, resulting in a tiny corresponding  $\sigma_{k,\beta}^*$ . In all these cases, we simply choose a fraction of the current  $\sigma_k$ . On the other hand, note that, by definition, the values of  $\sigma_{k,\beta}^*$  computed in (3.24) and (3.25) are positive and smaller than the current  $\sigma_k$ . Figure 3.4 shows the value of  $\sigma_{k+1}$  computed by Algorithm 3.1 as a function of the objective function value  $f(x_k + s_k)$ . This curve for  $\sigma_{k+1}$  is a piecewise linear function where the sloping pieces correspond to values of  $\sigma_{k+1}$  computed by the interpolation rules (3.24)–(3.26).

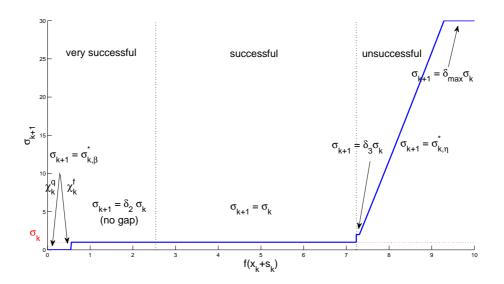


Figure 3.4: Plot of  $\sigma_{k+1}$ , computed by Algorithm 3.1 with parameters  $\beta = 0.01$ ,  $\alpha_{max} = 2$ ,  $\epsilon_{\chi} = 10^{-8}$ ,  $\delta_1 = 0.1$ ,  $\delta_2 = 1$ ,  $\eta = \eta_1$ ,  $\delta_3 = 2$ ,  $\delta_{max} = 30$ , as a function of  $f(x_k + s_k)$ .

## 4 Numerical experiments

We now present numerical experiments on nonlinear least-squares problems (1.2), where we study the numerical behaviour of the trust-region and the ARC algorithms employing the different updating rules presented in Section 3 in a first stage, and, in a second stage, compare the two algorithms using the best performing rules.

To compare the overall computational effort of the algorithms we use the performance profiles proposed by Dolan and Moré [9] for a given set of test problems and a given selection of algorithms. For each problem P in our testing set and each Algorithm A, we let  $fe_{P,A}$  denote the number of function evaluations required to solve problem P using Algorithm A and  $fe_P$  be number of function evaluations required by the best algorithm to solve problem P, i.e. the algorithm which uses the fewest function evaluations. The performance profile is defined for the algorithm A as

$$\pi_A(\tau) = \frac{\text{number of problems s.t. } \mathbf{fe}_{P,A} \le \tau \, \mathbf{fe}_P}{\text{number of problems}}, \quad \tau \ge 1.$$
(4.1)

In what follows and in order to improve readability of the performance profile graphs, we limit the plot  $\pi_A(\tau)$  to the interval [1, 4] and report the number of failures in the legend.

#### 4.1 The problem set

Numerical results are given for problems from the CUTEr test collection [15]. The test examples we consider are constructed using the CUTEr interactive select tool in order to locate the problems with no objective function and with constraints that are systems of nonlinear equations. We exclude the problems CHEMRCTA, CHEMRCTB, DRCAVTY3,

#### Algorithm 3.1: Regularization parameter update

Given the current  $x_k, s_k, \sigma_k$ , let the constants  $\eta_1$  and  $\eta_2$  be fixed by Algorithm 2.1. Let the positive threshold  $\epsilon_{\chi}$  and the constants  $\delta_1, \delta_2, \delta_3, \delta_{max}, \beta, \eta$  be chosen such that

$$0 < \delta_1 < \delta_2 \le 1 \le \delta_3 \ll \delta_{max}, \ 0 < \beta < 1, \ 0 < \eta < 3/2, \ 1 < \alpha_{max}$$

Compute  $\rho_c(s_k)$  by (2.7) and

$$\chi_k = c_k(s_k) - \max\{f(x_k + s_k), q_k(s_k)\}$$

• If 
$$\rho_c(s_k) \ge 1$$
 and  $\chi_k \ge \epsilon_{\chi}$ , then

- If  $f(x_k + s_k) \ge q_k(s_k)$ , solve equation (3.14) with  $\chi_k^f = \chi_k$ . Let  $\mathcal{A}^* = \{ \alpha \mid \alpha \text{ is a root of } (3.14) \text{ and } \alpha \ge \sqrt[3]{\beta} \}.$ 
  - \* If  $\mathcal{A}^* = \emptyset$ , set  $\sigma_{k+1} = \max\{\delta_1 \sigma_k, \epsilon_m\}.$
  - \* If  $\mathcal{A}^* \neq \emptyset$ , let  $\alpha_{\beta}^* = \operatorname{argmin}\{(\alpha \sqrt[3]{\beta}) \mid \alpha \in \mathcal{A}^*\}$ . If  $\alpha_{\beta}^* \leq \alpha_{max}$ , compute

$$\sigma_{k,\beta}^* = \sigma_k + 3 \frac{\chi_k}{\|s_k\|^3} \left( \frac{\beta - \alpha_\beta^{*3}}{\alpha_\beta^{*3}} \right), \qquad (3.24)$$

and set  $\sigma_{k+1} = \max\{\sigma_{k,\beta}^*, \epsilon_m\};$ If  $\alpha_{\beta}^* > \alpha_{max}$ , set  $\sigma_{k+1} = \max\{\delta_1 \sigma_k, \epsilon_m\}.$ 

- Else if  $f(x_k + s_k) < q_k(s_k)$ , solve equation (3.19) with  $\chi_k^q = \chi_k$ . Let  $\mathcal{A}^* = \{ \alpha \mid \alpha \text{ is a root of (3.19) and } \alpha \geq \sqrt[3]{\beta} \}.$ 

- \* If  $\mathcal{A}^* = \emptyset$ , set  $\sigma_{k+1} = \max\{\delta_1 \sigma_k, \epsilon_m\}.$
- \* If  $\mathcal{A}^* \neq \emptyset$ , let  $\alpha_{\beta}^* = \operatorname{argmin}\{(\alpha \sqrt[3]{\beta}) \mid \alpha \in \mathcal{A}^*\}$ . If  $\alpha_{\beta}^* \leq \alpha_{max}$ , compute

$$\sigma_{k,\beta}^* = \frac{\beta}{\alpha_\beta^{*3}} \sigma_k, \qquad (3.25)$$

and set  $\sigma_{k+1} = \max\{\sigma_{k,\beta}^*, \epsilon_m\};$ If  $\alpha_{\beta}^* > \alpha_{max}$ , set  $\sigma_{k+1} = \max\{\delta_1 \sigma_k, \epsilon_m\}.$ 

- Else if  $\rho_c(s_k) \ge 1$  and  $\chi_k < \epsilon_{\chi}$ , set  $\sigma_{k+1} = \max\{\delta_2 \sigma_k, \epsilon_m\}$ .
- Else if  $\rho_c(s_k) \in [\eta_2, 1)$ , set  $\sigma_{k+1} = \max\{\delta_2 \sigma_k, \epsilon_m\}$ .
- Else if  $\rho_c(s_k) \in [\eta_1, \eta_2)$ , set  $\sigma_{k+1} = \sigma_k$ .
- Else if  $\rho_c(s_k) \in [0, \eta_1)$ , set  $\sigma_{k+1} = \delta_3 \sigma_k$ .
- Else  $(\rho_c(s_k) < 0)$ , compute the positive root  $\alpha_n^*$  of equation (3.23) and compute

$$\sigma_{k,\eta}^* = \frac{-g_k^T s_k - s_k^T H_k s_k \alpha_\eta^*}{\alpha_\eta^{*2} \|s_k\|^3}.$$
(3.26)

Set  $\sigma_{k+1} = \min\{\max\{\sigma_{k,\eta}^*, \delta_3\sigma_k\}, \delta_{max}\sigma_k\}.$ 

FLOSP2HH, FLOSP2HL, FLOSP2HM, FLOSP2TH, FLOSP2TL, FLOSP2TM, HYDCAR2O, SEMICON2 and SEMICN2U as no algorithm succeeded in solving these problems for any tested parameter choice. For some CUTEr problems, we considered variants that differ in the dimensions (denoted with the superscript  $^{2}$ ,<sup>3</sup>). The resulting testing set consists of 95 problems of the form (1.2) whose names and dimensions are reported in Table 5.1 of Appendix. The problems ARGLALE, ARGBLE, GROWTH, HIMMELBD and OSCIPANE are large residual problems, i.e. the objective function value at the computed solution is much greater than one, the remaining are small or zero residual problems. Moreover, for 9 problems m > n, for 28 problems m < n, the remaining 58 ones being square.

#### 4.2 Implementation issues

We implemented Algorithm 2.1 in Fortran 95, using the procedures presented in Section 2.1 to solve the subproblem at Step 1. We consider two implementations of the trust-region algorithm (TR-ST and TR-bST) which use the GALAHAD's package [16] LSTR and differ in the computation of the boundary trust-region solution: TR-ST computes the Steihaug-Toint point, TR-bST computes a more accurate solution as described in Section 2.1. The tested version of the ARC algorithm for solving problem (1.2) has been implemented using the GALAHAD's packages LSRT and it is denoted by ARC-LS.

In Algorithm 2.1, we set the specific algorithmic constants

$$\eta_1 = 0.01, \ \eta_2 = 0.95, \tag{4.2}$$

and the initial regularization parameters  $\Delta_0$  and  $\sigma_0$  are chosen equal to one. The algorithm is terminated as soon as either

$$||J_k^T h_k|| \le \max\{\epsilon_{ga}, \epsilon_{gr} ||J_0^T h_0||\}$$
 or  $||h_k|| \le \max\{\epsilon_{fa}, \epsilon_{fr} ||h_0||\},$  (4.3)

where  $\epsilon_{fa}, \epsilon_{ga}, \epsilon_{fr}, \epsilon_{gr} > 0$  are tolerances chosen as  $\epsilon_{fa} = \epsilon_{ga} = 10^{-6}$ ,  $\epsilon_{fr} = \epsilon_{gr} = 10^{-12}$ . Moreover, we require that the trial step  $s_k$  computed at Step 1 of Algorithm 2.1 satisfies the inexact stopping criterion given by

$$\|\nabla m_k(s_k)\| \le \min\{\epsilon_{in}, \|\nabla m_k(0)\|^{1/2}\} \|\nabla m_k(0)\|,$$
(4.4)

where  $m_k$  represents the models  $c_k$  in (2.11) and  $q_k$  in (2.10) and  $\epsilon_{in} = 10^{-1}$  is fixed. If the problem dimension n is lower than 50, we allow for the generation of the full space in the Krylov sequence in order to compute a very accurate solution of the subproblems (2.12) and (2.13). Furthermore, any run exceeding 2 hours of CPU time, performing more than 5000 outer iterations or if the magnitude of computed search direction is lower than  $10\epsilon_m$ , is considered a failure. All other parameters in the GALAHAD's packages are set at their default values.

All our tests were performed on an Intel Xeon (TM) 3.4 Ghz, 1GB of RAM; the codes are all double precision, and compiled under g95 without optimization (default).

#### 4.3 Numerical results

We consider first the trust-region algorithm and the trust-region radius updating rules described in Section 3. In particular we compare the updating rules (3.1) and (3.4) where we used parameter values given by

$$\gamma_1 = 1/2, \ \gamma_2 = 2, \ \gamma_3 = 0.0625,$$
(4.5)

and tried the values  $\eta_1$  and  $\eta_2$  given in (4.2) for the parameter  $\eta$  in (3.3).

In Figure 4.1, the function evaluation performance profiles show that both TR-ST and TR-bST are slightly more efficient using the updating rule (3.4) with  $\eta = \eta_1$ . Moreover, TR-bST is also a little more robust with this choice. The performance profile of Figure 4.2 summarizes the comparison between the two trust-region implementations using the best performing rule with the best parameter choice. As one might hope, the figure suggests that the extra effort required to solve the subproblem more accurately appears to offer some overall benefit.

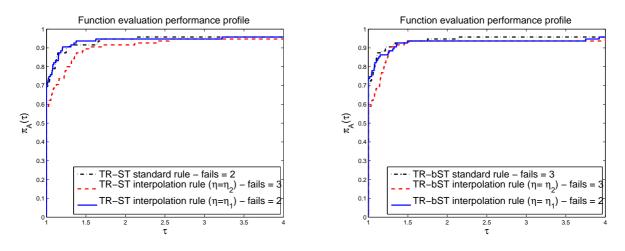


Figure 4.1: The function evaluation performance profile: TR-ST (left) and TR-bST (right) with (3.1) ("standard rule") and (3.4) using  $\eta = \eta_1, \eta_2$  ("interpolation rule").

We now examine the sensitivity in number of function evaluations for the parameter choices of the new updating rule for  $\sigma_k$  for the ARC algorithm. To this purpose, we performed a small parametric study starting from the following reasonable values for the parameters in Algorithm 3.1

$$\beta = 1/100, \ \alpha_{max} = 2, \ \epsilon_{\chi} = 10^{-8}, \ \delta_1 = 1/10, \ \delta_2 = 1, \ \eta = \eta_1, \ \delta_3 = 2, \ \delta_{max} = 100, \ (4.6)$$

and varying one parameter at the time in some set to find the best performing value.

More precisely, let all the parameters be ordered as  $\beta$ ,  $\alpha_{max}$ ,  $\epsilon_{\chi}$ ,  $\delta_1$ ,  $\delta_2$ ,  $\eta$ ,  $\delta_3$ ,  $\delta_{max}$  and be fixed as in (4.6). Let p be a parameter to be analyzed. Moreover, let  $I_p = \{p_1, \ldots, p_q\}$ be a set of trial values for p,  $A_{p_i}$  be the ARC-LS algorithm run with  $p = p_i$  and let  $\pi_{A_{p_i}}(\tau)$ be the performance measure defined in (4.1) comparing the algorithms  $A_{p_i}$ ,  $p_i \in I_p$ . To

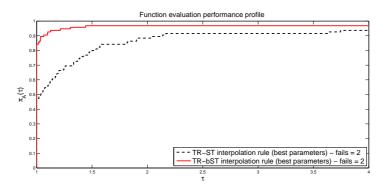


Figure 4.2: The function evaluation performance profile: TR-ST and TR-bST with the interpolation rule (3.4) and the best parameter choice ( $\eta = \eta_1$ ).

estimate the efficiency of these algorithms, we compute the percentage of problems  $(\%pb_{\hat{\tau}})$ for which  $\pi_{A_{p_i}}(\tau) \leq \hat{\pi}$  with  $\hat{\pi} \gtrsim 1$  and to evaluate their robustness, we compute the number of failures. Taking into account these performance measures, we fix the "best" value for the parameter  $p \in I_p$  and we proceed with the analysis of the subsequent parameter in the list. In Table 4.1, we report the sets  $I_p$  for all the parameters in Algorithm 3.1, the efficiency measure  $(\%pb_{\hat{\tau}})$  for  $\hat{\tau} = 1, 1.15, 1.25, 1.5, 2$  and the number of failures (**#fails**). We note that a more sophisticated choice, in which the globally optimal parameters for our test set is determined [1], is possible but has not been performed.

For each set  $I_p$ , it is quite easy to find the best performing parameter choice. It results from Table 4.1 that the new updating rule is not very sensitive to the parameter choice and that ARC-LS performs slightly better with the following parameter assignment:

$$\beta = 1/100, \ \alpha_{max} = 2, \ \epsilon_{\chi} = 10^{-10}, \ \delta_1 = 1/10, \ \delta_2 = 1, \ \eta = \eta_1, \ \delta_3 = 2, \ \delta_{max} = 100.$$
 (4.7)

We remark that in the experiments, a solution  $\alpha_{\beta}^*$  of equations (3.14) and (3.19) was always found and that only in few cases this values was larger than  $\alpha_{max}$ . Moreover, the value  $\sigma_{k,\eta}^*$ computed by (3.26) was very often positive and lower than the current  $\sigma_k$ . Consequently, the regularization parameter was in fact updated by using the proposed interpolation techniques most of the time.

In Figure 4.3, ARC-LS using Algorithm 3.1 and the parameters in (4.7) is compared with ARC-LS using the old rule (3.5) and  $\gamma = 2$  employed in [3]. The new rule clearly outperforms the old one. A possible explanation of the relatively poor behaviour of ARC-LS with the old rule may be found in what follows. In the experiments, we noticed that the norm of the gradient oscillates considerably for some problems, resulting in high oscillations in the updated  $\sigma_k$  through the iterations. Furthermore, we observed that, using (3.5),  $\sigma_k$  was updated in several runs using a small  $||g_k||$  and hence was considerably reduced; the next iterate was then unsuccessful and doubling  $\sigma_k$  to recover an acceptable  $\sigma_k$  gave rise to many unsuccessful iterations.

p	$I_p$	#fails	$\% pb_{\hat{\tau}}, \hat{\tau} = 1$	$\% pb_{\hat{\tau}}, \hat{\tau} = 1.15$	$\% pb_{\hat{\tau}}, \hat{\tau} = 1.25$	$\% pb_{\hat{\tau}}, \hat{\tau} = 1.5$	$\% pb_{\hat{\tau}}, \hat{\tau} = 2$
	0.001	3	58.95	90.53	92.63	93.68	95.79
	0.005	4	55.79	86.32	92.63	93.68	95.79
β	0.01	3	68.42	91.58	95.79	95.79	96.84
	0.05	3	51.58	82.11	89.47	93.68	95.79
	0.1	4	51.58	81.05	89.47	95.79	95.79
	1	4	42.11	69.47	75.79	91.58	94.74
	2	3	60.00	88.42	92.63	94.74	94.74
	3.5	3	55.79	85.26	90.53	92.63	92.63
$\alpha_{max}$	5	3	61.05	84.21	90.53	92.63	92.63
	10	3	63.16	86.32	90.53	91.58	92.63
	50	4	61.05	84.21	89.47	90.53	91.58
	$10^{-12}$	3	81.05	92.63	93.68	95.79	95.79
	$10^{-11}$	3	83.16	93.68	95.79	95.79	96.84
	$10^{-10}$	2	80.00	95.79	96.84	96.84	97.89
$\epsilon_{\chi}$	$10^{-9}$	3	76.84	92.63	93.68	94.74	96.84
	$10^{-8}$	3	73.68	92.63	92.63	94.74	94.74
	$10^{-6}$	4	65.26	78.95	82.11	86.32	90.53
	0.01	6	69.47	85.26	88.42	91.58	92.63
	0.05	3	65.26	92.63	94.74	95.79	95.79
$\delta_1$	0.1	2	71.58	94.74	97.89	97.89	97.89
	0.25	3	62.11	87.37	93.68	95.79	95.79
	0.5	3	58.95	83.16	92.63	95.79	95.79
	0.25	3	61.05	74.74	85.26	91.58	94.74
	0.5	3	53.68	78.95	85.26	90.53	96.84
$\delta_2$	0.75	4	57.89	86.32	91.58	94.74	95.79
	0.9	4	57.89	85.26	90.53	93.68	95.79
	1	2	58.95	89.47	95.79	97.89	97.89
	$\eta_1$	2	72.63	94.74	95.79	96.84	97.89
	$(\eta_2 - \eta_1)/2$	4	68.42	88.42	91.58	94.74	95.79
$\eta$	$\eta_2$	5	62.11	81.05	88.42	91.58	94.74
	1.25	3	63.16	78.95	87.37	89.47	90.53
	1.50	4	69.47	89.47	93.68	94.74	94.74
	2	2	72.63	93.68	96.84	97.89	97.89
$\delta_3$	2.5	4	66.32	89.47	91.58	94.74	95.79
	3	4	65.26	88.42	94.74	95.79	95.79
	4	3	66.32	83.16	92.63	94.74	96.84
	10	4	66.32	86.32	90.53	91.58	93.68
	50	4	70.53	89.47	93.68	95.79	95.79
$\delta_{max}$	100	2	74.74	95.79	97.89	97.89	97.89
	500	4	71.58	91.58	93.68	93.68	94.74
	1000	4	72.63	88.42	93.68	93.68	94.74

Table 4.1: Parametric study.

Finally, we compare TR-ST, TR-bST and ARC-LS using the best performing updating rules for the regularization parameters, i.e. for the trust-region radius  $\Delta_k$  the rule (3.4) with the parameters in (4.5) and  $\eta = \eta_1$  and for the regularization parameter  $\sigma_k$ , the rule presented in Algorithm 3.1 with the parameter choice (4.7). The corresponding function evaluation performance profiles are plotted in Figures 4.4 and 4.5. ARC-LS fails on problems ARWHDNE, DRCAVITY2, TR-bST on problems DRCAVITY2, POROUS2 and TR-ST on problems QR3D<sup>2</sup>, POROUS2. Evidently, ARC-LS is much more efficient than TR-ST. Compared to TR-bST, it is better 68.42% of the runs and TR-bST is within a factor 2 of ARC-LS for the 88.10% of the runs.

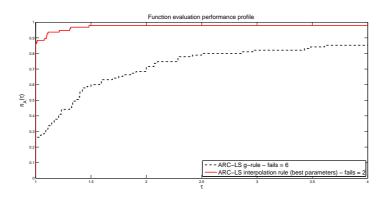


Figure 4.3: The function evaluation performance profile: ARC-LS with (3.5) ("g-rule") and ARC-LS with Algorithm 3.1 and parameters (4.7) ("interpolation rule (best parameters)").

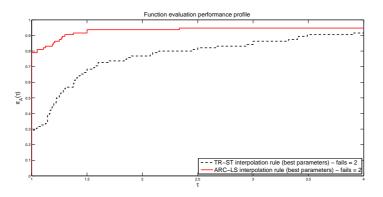


Figure 4.4: The function evaluation performance profile: TR-ST rule (3.4) with  $\eta = \eta_1$  ("interpolation rule (best parameters)") and ARC-LS with Algorithm 3.1 and parameters (4.7) ("interpolation rule (best parameters)").

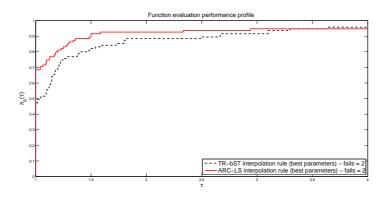


Figure 4.5: The function evaluation performance profile: TR-bST rule (3.4) with  $\eta = \eta_1$  ("interpolation rule (best parameters)") and ARC-LS with Algorithm 3.1 and parameters (4.7) ("interpolation rule (best parameters)").

We report in the Appendix the complete set of results of the experiments described in this section.

We also considered strategies for choosing the initial regularization parameter  $\sigma_0$  along the lines of the strategy proposed in [22] for automatically computing the initial trustregion radius. In particular, we tested a strategy in which one solves a one-dimensional minimization problem (along the steepest descent direction) in the hope of estimating a better value of  $\sigma_0$  for starting the minimization in the full space. However, these experiments (not reported here) produced disappointing results in that it turned out to be generally better to start minimization in the full-space from the start and not "waste" additional function evaluations for estimating  $\sigma_0$ . This is not entirely unexpected in our context where we assume the cost of function evaluation to dominate the inner linear algebra calculations. But it is also clear that any *a priori* user estimation of the Hessian Lipschitz constant can be usefully exploited by selecting  $\sigma_0$  appropriately.

## 5 Conclusion

In this paper we propose a new reliable strategy to update the regularization parameter in the cubic regularization algorithm (ARC). This strategy is based on analyzing the adequacy between the objective function and its cubic model, and exploits its overestimation property. Moreover, it has the favorable feature of not requiring extra function values. We report numerical tests which show that the new rule considerably improves the numerical performance of the ARC algorithm. We also provide a numerical comparison between the ARC and trust-region frameworks on a set of large nonlinear least-squares CUTEr problems. These suggest a numerical advantage of the former on our set of test problems.

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## Appendix

Table 5.1 contains the problem set information (name and dimensions). Tables 5.2-5.6 collect all the results of the experiments described in Section 4: we reported the total number of function evaluation for each method and algorithmic option tested and we used the following symbols for the failures: '\*' for the time exceeding runs, '>' for the runs exceeding the maximum number of iteration allowed, 'ss' if the norm of the search step is below the fixed treshold.

Name	n	m	Name	n	m	Name	n	m
AIRCRFTA	8	5	DECONVNE	61	41	OSCIPANE	500	500
ARGAUSS	3	15	DRCAVTY1	196	100	PFIT1	2	2
ARGLALE	200	400	DRCAVTY2	4489	3969	PFIT2	2	2
ARGLBLE	200	400	EIGENA	110	110	PFIT3	2	2
ARGTRIG	200	200	EIGENA <sup>2</sup>	2550	2550	PFIT4	2	2
ARTIF	502	500	EIGENA <sup>3</sup>	4970	4970	POROUS1	1024	900
$ARTIF^2$	5002	5000	EIGENB	110	110	$POROUS1^2$	5184	4900
ARWDHNE	500	998	$EIGENB^2$	2550	2550	POROUS1 <sup>3</sup>	22500	21904
BDVALUES	102	100	EIGENC	462	462	POROUS2	1024	900
$BDVALUS^2$	5002	5000	$EIGENC^2$	2652	2652	$POROUS2^2$	5184	4900
BOOTH	2	2	GOTTFR	2	2	$POROUS2^3$	22500	21904
BRATU2D	484	400	GROWTH	3	12	$POROUS2^4$	62500	61504
$BRATU2D^2$	5184	4900	HATFLDF	3	3	POWELLBS	2	2
BRATU2DT	484	400	HATFLDG	25	25	POWELLSQ	2	2
$BRATU2DT^2$	5184	4900	HEART6	6	6	QR3D	610	610
BRATU3D	1000	512	HEART8	8	8	$QR3D^2$	2420	2420
$BRATU3D^2$	4913	3375	HIMMELBA	2	2	QR3DBD	457	610
BROWNALE	200	200	HIMMELBC	2	2	$QR3DBD^2$	1717	2420
$BROWNALE^2$	1000	1000	HIMMELBD	2	2	RECIPE	3	3
BROYDN3D	1000	1000	HIMMELBE	3	3	SINVALNE	2	2
$BROYDN3D^2$	10000	10000	HS8	2	2	SPMSQRT	10000	16664
BROYDNBD	1000	1000	HYDCAR6	29	29	TRIGGER	7	6
$BROYDNBD^2$	10000	10000	HYPCIR	2	2	WOODSNE	10000	7501
CBRATU2D	3200	2888	INTEGREQ	102	100	YATP1SQ	2600	2600
CBRATU3D	3456	2000	$INTEGREQ^2$	502	500	$YATP1SQ^2$	40400	40400
CHANDHEQ	100	100	METHANB8	31	31	$YATP1SQ^{3}$	63000	63000
CHANNEL	2400	2398	METHANL8	31	31	YATP2SQ	2600	2600
$CHANNEL^2$	9600	9598	MSQRTA	4900	4900	$YATP2SQ^2$	40400	40400
CHNRSBNE	50	98	$MSQRTA^2$	5625	5625	$YATP2SQ^3$	63000	63000
CLUSTER	2	2	MSQRTB	4900	4900	YFITNE	3	17
COOLHANS	9	9	$MSQRTB^2$	5625	5625	ZANGWIL3	3	3
CUBENE	2	2	NYSTROM5	18	20			

Table 5.1: The problem set.

ARTUP <sup>2</sup> 25         25         25         17         14         17         14         14         <		TR-ST			Т	R-bST				TR-ST		Т	R-bST	
AIRCRETA         4<	Name	standard	interp	olation	standard	inter	polation	Name	standard	interp	olation	standard	interp	olation
ARGALSS       2       2       2       2       2       2       2       1       I       1       2       5       6 </td <td></td> <td></td> <td><math>\eta_2</math></td> <td><math>\eta_1</math></td> <td></td> <td><math>\eta_2</math></td> <td><math>\eta_1</math></td> <td></td> <td></td> <td><math>\eta_2</math></td> <td><math>\eta_1</math></td> <td></td> <td><math>\eta_2</math></td> <td><math>\eta_1</math></td>			$\eta_2$	$\eta_1$		$\eta_2$	$\eta_1$			$\eta_2$	$\eta_1$		$\eta_2$	$\eta_1$
ARCLALE       6       6       6       6       6       6       6       111MMELBD       41       4			4	4		4	4	HIMMELBA	4	4	4	4	4	4
ARGUBLE       5       5       5       5       5       5       6       HIMMELBE       4		2	2	2	2	2	2	HIMMELBC	6	6	6	6	6	
ARCTIF       9       60       425       33       34       34       34       34       34       34       34       34       34       34       34       34       34       34       34       34       34       34       34 <th< td=""><td>ARGLALE</td><td>6</td><td>6</td><td>6</td><td>6</td><td>6</td><td>6</td><td>HIMMELBD</td><td>41</td><td>25</td><td>30</td><td>42</td><td>24</td><td>30</td></th<>	ARGLALE	6	6	6	6	6	6	HIMMELBD	41	25	30	42	24	30
ARTIFF         20         20         22         20         22         11YDCAR6         458         530         600         425         381         438           ARTIFF         25         25         25         25         25         17         17         17         IYPCIR         5         6         7         7	ARGLBLE	5	5	5	5	5	5	HIMMELBE	4	4	4	4	4	4
ARTUP <sup>2</sup> 25         25         25         17         15         5	ARGTRIG	9	9	9	9	9	9	HS8	7	7	7	6	6	6
$\begin{array}{c c c c c c c c c c c c c c c c c c c $		20	20	20	22	20	22	HYDCAR6	458	530	600	425	381	438
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	$ARTIF^2$	25	25	25	17	17	17	HYPCIR	5	5	5	5	5	5
BDVALUES <sup>2</sup> 391         391         416         416         416         416         416         416         416         416         416         416         416         416         416         44 <th< td=""><td>ARWDHNE</td><td>398</td><td>SS</td><td>SS</td><td>ss</td><td>ss</td><td>172</td><td>INTEGREQ</td><td>5</td><td>5</td><td>5</td><td>5</td><td>5</td><td>5</td></th<>	ARWDHNE	398	SS	SS	ss	ss	172	INTEGREQ	5	5	5	5	5	5
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	BDVALUES	43	43	43	62	62	62	$INTEGREQ^2$	5	5	5	5	5	5
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	$BDVALUES^2$	391	391	391	416	416	416	METHANB8	94	94	94	94	94	94
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	BOOTH	4	4	4	4	4	4	METHANL8	237	189	204	266	178	266
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	BRATU2D	7	7	7	5	5	5		46	44	45	44	46	46
BRATU2D <sup>2</sup> 9       9       9       6       6       6       MSQRTB <sup>2</sup> 52       51       50       47       54       49         BRATU3D       8       8       8       7       7       7       NYSTROM5       223       268       198       140       149       129         BRATU3D <sup>2</sup> 10       10       10       8       8       8       CECPANE       8       11       11       9       10       10 </td <td><math>BRATU2DT^2</math></td> <td>12</td> <td>12</td> <td>12</td> <td>10</td> <td>10</td> <td>10</td> <td></td> <td>54</td> <td>57</td> <td>55</td> <td>54</td> <td>61</td> <td>53</td>	$BRATU2DT^2$	12	12	12	10	10	10		54	57	55	54	61	53
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	BRATU2DT	24	19	22	14	13	14	MSQRTB	47	44	47	44	44	44
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	$BRATU2D^2$	9	9	9	6	6	6	$MSQRTB^2$	52	51	50	47	54	49
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	BRATU3D	8	8	8	7	7	7	NYSTROM5	223	268	198	140	149	129
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	$BRATU3D^2$	10	10	10	8	8	8	OSCIPANE	8	8	8	8	8	8
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	BROWNALE	6	6	6	6	6	6	PFIT1	13	105	109	13	105	51
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	$BROWNALE^2$	8	8	8	8	8	8	PFIT2	28	13	13	28	13	13
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	BROYDN3D	9	9	9	9	9	9	PFIT3	10	11	9	11	11	9
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	BROYDN3D <sup>2</sup>	11	11	11	11	11	11	PFIT4	14	169	14	9	9	9
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	BROYDNBD	19	19	19	18	18	18	POROUS1	47	51	44	41	39	35
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	$BROYDNBD^2$	27	29	27	19	19	19	POROUS1 <sup>2</sup>	190	251	147	90	86	81
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	CBRATU2D	8	8	8	6	6	6	POROUS1 <sup>3</sup>	822	1016	888	168	183	154
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	CBRATU3D	9	9	9	8	8	8	POROUS2	2174	149	1788	>	>	>
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	CHANDHEQ	15	15	15	14	14	14	$POROUS2^2$	195	291	230	106	137	120
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	CHANNEL	213	451	293	154	185	154	POROUS2 <sup>3</sup>	1045	1204	914	200	281	186
CLUSTER8888889090191810919108COOLHANS890827538604777653QR3D621497573186165153CUBENE66666QR3D <sup>2</sup> >>>>915700916DECONVNE161816121212QR3DBD342307353749274DRCAVTY1414441354035QR3DBD <sup>2</sup> 125212341249678506686DRCAVTY2**490***RECIPE243339182118EIGENA2121202020SINVALNE2730242530735333EIGENA <sup>2</sup> 108109106727972SPMSQRT151216 </td <td><math>CHANNEL^2</math></td> <td>270</td> <td>381</td> <td>159</td> <td>106</td> <td>103</td> <td>115</td> <td><math>POROUS2^4</math></td> <td>2749</td> <td>3124</td> <td>2478</td> <td>299</td> <td>364</td> <td>271</td>	$CHANNEL^2$	270	381	159	106	103	115	$POROUS2^4$	2749	3124	2478	299	364	271
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	CHNRSBNE	61	75	64	48	60	48	POWELLBS	87	113	82	69	79	74
CUBENE666666 $(QR3D^2)$ >>>>915700916DECONVNE161816121212QR3DBD342307353749274DRCAVTY1414441354035QR3DBD21252123412496678506686DRCAVTY2**400****RECIPE243339182118EIGENA212121202020SINVALNE273024253024EIGENA22108109106727972SPMSQRT151515151515EIGENA3166171171848784TRIGGER888888EIGENB131176141133154143WOODSNE423939353333EIGENC293102104707966YATP1SQ2282930262725EIGENC293102104707966YATP1SQ3282727262425GOTTFR111611111611YATP2SQ2303030313131GROWTH547154117154YATP2SQ	CLUSTER	8	8	8	8	8	8	POWELLSQ	98	19	18	109	19	108
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	COOLHANS	890	827	538	604	777	653	QR3D	621	497	573	186	165	153
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	CUBENE	6	6	6	6	6	6	$QR3D^2$	>	>	>	915	700	916
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	DECONVNE	16	18	16	12	12	12	QR3DBD	342	307	353	74	92	74
EIGENVERT2       12       21       21       21       21       20       20       20       SINVALNE       27       30       24       25       30       24       30       33       33       33       33       33       33       33       33       33       33       33       33       33       33<	DRCAVTY1	41	44	41	35	40	35	$QR3DBD^2$	1252	1234	1249	678	506	686
EIGENA2108109106727972SPMSQRT151515151515EIGENA3166171171848784TRIGGER888888EIGENB131176141133154143WOODSNE423939353333EIGENB2104712831050687862936YATP1SQ414736293029EIGENC93102104707966YATP1SQ2282930262725EIGENC2834751904245263249YATP1SQ3282727262425GOTTFR111611116111YATP2SQ2303030313131GROWTH547154117154YATP2SQ3303030313131HATFLDF9232983230YATP2SQ3303030313131HATFLDG8889119YFITNE465062465061HEART6484558528617687580ZANGWIL388888888	DRCAVTY2	*	*	490	*	*	*	RECIPE	24	33	39	18	21	18
EIGENA3166171171848784TRIGER8888888EIGENB131176141133154143WOODSNE423939353333EIGENB2104712831050687862936YATP1SQ414736293029EIGENC93102104707966YATP1SQ2282930262725EIGENC2834751904245263249YATP1SQ3282727262425GOTTFR111611111611YATP2SQ303030313131GROWTH547154117154YATP2SQ3303030313131HATFLDF9232983230YATP2SQ3303030313131HATFLDG8889119YFITNE465062465061HEART6484558528617687580ZANGWIL388888888	EIGENA	21	21	21	20	20	20	SINVALNE	27	30	24	25	30	24
EIGENB131176141133154143WOODSNE423939353333EIGENB2104712831050687862936YATP1SQ414736293029EIGENC93102104707966YATP1SQ2282930262725EIGENC2834751904245263249YATP1SQ3282727262425GOTTFR111611111611YATP2SQ303030313131GROWTH547154117154YATP2SQ2343434333333HATFLDF9232983230YATP2SQ3303030313131HATFLDG8889119YFITNE465062465061HEART6484558528617687580ZANGWIL38888888	EIGENA <sup>2</sup>	108	109	106	72	79	72	SPMSQRT	15	15	15	15	15	15
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	EIGENA <sup>3</sup>	166	171	171	84	87	84	TRIGGER	8	8	8	8	8	8
EIGENB2104712831050687862936YATP1SQ414736293029EIGENC93102104707966YATP1SQ2282930262725EIGENC2834751904245263249YATP1SQ3282727262425GOTTFR111611111611YATP2SQ3030303131GROWTH547154117154YATP2SQ2343434333333HATFLDF9232983230YATP2SQ3303030313131HATFLDG8889119YFITNE465062465061HEART6484558528617687580ZANGWIL38888888	EIGENB	131	176	141	133	154	143	WOODSNE	42	39	39	35	33	33
EIGENC93102104707966YATP1SQ2282930262725EIGENC2834751904245263249YATP1SQ3282727262425GOTTFR111611111611YATP2SQ303030313131GROWTH547154117154YATP2SQ2343434333333HATFLDF9232983230YATP2SQ3303030313131HATFLDG8889119YFITNE465062465061HEART6484558528617687580ZANGWIL38888888			1283	1050	687	862	936		41	47	36	29		
EIGENC <sup>2</sup> 834         751         904         245         263         249         YATP1SQ <sup>3</sup> 28         27         27         26         24         25           GOTTFR         11         16         11         11         16         11         YATP2SQ         30         30         30         31         31         31           GROWTH         54         71         54         11         71         54         YATP2SQ <sup>2</sup> 34         34         34         33         33         33           HATFLDF         9         23         29         8         32         30         YATP2SQ <sup>3</sup> 30         30         30         31         31         31           HATFLDF         9         23         29         8         32         30         YATP2SQ <sup>3</sup> 30         30         30         31         31         31           HATFLDG         8         8         8         9         11         9         YFITNE         46         50         62         46         50         61           HEART6         484         558         528         617         687         580         ZANGWIL3 <td>EIGENC</td> <td></td> <td></td> <td></td> <td></td> <td></td> <td></td> <td></td> <td>28</td> <td>29</td> <td>30</td> <td>26</td> <td>27</td> <td>25</td>	EIGENC								28	29	30	26	27	25
GOTTFR1116111611YATP2SQ303030313131GROWTH547154117154YATP2SQ <sup>2</sup> 343434333333HATFLDF9232983230YATP2SQ <sup>3</sup> 303030313131HATFLDG8889119YFITNE465062465061HEART6484558528617687580ZANGWIL3888888														
GROWTH547154117154YATP2SQ2343434333333HATFLDF9232983230YATP2SQ330303030313131HATFLDG8889119YFITNE465062465061HEART6484558528617687580ZANGWIL3888888														
HATFLDF       9       23       29       8       32       30       YATP2SQ <sup>3</sup> 30       30       30       31       31       31         HATFLDG       8       8       8       9       11       9       YFITNE       46       50       62       46       50       61         HEART6       484       558       528       617       687       580       ZANGWIL3       8       8       8       8       8       8       8														
HATFLDG       8       8       9       11       9       YFITNE       46       50       62       46       50       61         HEART6       484       558       528       617       687       580       ZANGWIL3       8 <td< td=""><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td></td<>														
HEART6         484         558         528         617         687         580         ZANGWIL3         8 <td>HATFLDG</td> <td></td>	HATFLDG													
HEAR18 40 49 55 58 40 48	HEART8	46	49	53	38	46	48							

			β			α <sub>max</sub>								$\delta_1$					
Name	0.001	0.005	0.01	0.5	0.1	1	3.5	5	10	50	$10^{-12}$	$10^{-11}$	$\epsilon_{\chi}$ $10^{-10}$	$10^{-9}$	$10^{-6}$	0.01	0.05	0.25	0.5
AIRCRFTA	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4
ARGAUSS	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3
ARGLALE	6	6	6	6	6	6	5	5	4	4	6	6	6	6	6	5	5	6	7
ARGLBLE	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
ARGTRIG	9	9	9	9	9	9	9	9	9	9	9	9	9	9	9	9	9	9	9
ARTIF	21	21	21	21	21	21	21	21	21	21	21	21	21	21	21	20	20	23	24
$ARTIF^2$	23	20	20	20	23	23	19	19	18	18	20	20	20	20	20	19	19	19	21
ARWDHNE	ss	SS	SS	SS	ss	SS	ss	ss	ss	ss	ss	SS	SS	SS	ss	ss	SS	ss	ss
BDVALUES	39	63	47	49	44	39	47	47	47	47	34	34	34	34	47	34	34	34	34
BDVALUES <sup>2</sup>	*	*	*	*	*	*	*	*	*	270	*	*	270	*	*	*	*	*	*
BOOTH	4	4	4	4	5	5	4	4	4	4	4	4	4	4	4	4	4	4	4
BRATU2D	6	6	6	7	7	8	6	6	6	6	6	6	6	6	6	6	6	7	7
$BRATU2D^2$	6	6	6	9	9	10	6	6	6	6	6	6	6	6	6	6	6	6	6
BRATU2DT	14	15	14	15	15	19	14	14	14	14	14	14	14	14	14	14	14	14	14
$BRATU2DT^2$	10	10	10	11	13	13	10	10	10	10	10	10	10	10	10	10	10	10	10
BRATU3D	7	7	7	7	7	8	7	7	7	7	7	7	7	7	7	7	7	7	7
$BRATU3D^2$	7	7	7	7	8	8	7	7	7	7	7	7	7	7	7	7	7	7	7
BROWNALE	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4
BROWNALE <sup>2</sup>	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3
BROYDN3D	8	8	8	8	8	7	8	8	8	8	8	8	8	8	8	8	8	8	8
$BROYDN3D^2$	9	9	9	9	9	8	9	9	9	9	9	9	9	9	9	9	9	9	9
BROYDNBD	10	10	10	10	10	9	10	10	10	10	10	10	10	10	10	10	10	10	10
BROYDNBD <sup>2</sup>	11	11	11	11	11	10	11	11	11	11	11	11	11	11	11	11	11	11	11
CBRATU2D	6	7	7	7	9	9	7	7	7	7	7	7	7	7	7	7	7	7	7
CBRATU3D	7	7	7	7	8	8	7	7	7	7	7	7	7	7	7	7	7	7	7
CHANDHEQ	19	19	19	19	19	15	19	19	19	19	19	19	19	19	20	17	18	19	19
CHANNEL	142	147	140	146	136	147	145	145	142	142	139	139	140	140	206	143	147	148	148
CHANNEL <sup>2</sup>	99	97	94	95	98	102	101	101	101	101	101	101	94	94	94	106	101	98	95
CHNRSBNE	40	40	41	41	41	40	41	41	41	41	41	41	41	41	41	41	41	41	41
CLUSTER	10	10	10	10	10	9	10	10	10	10	10	10	10	10	11	10	10	10	10
COOLHANS	381	462	467	530	544	560	467	21	21	21	431	431	441	467	>	577	434	455	571
CUBENE	14	14	14	14	14	14	14	14	14	14	14	14	14	14	14	14	14	14	14
DECONVNE	38	36	36	38	39	13	36	36	36	36	17	17	19	33	36	19	19	19	19
DRCAVTY1	39	38	39	37	38	42	39	39	39	39	43	43	43	43	36	43	43	43	43
DRCAVTY2	*	*	*	*	*	*	*	*	*	*	*	*	*	*	*	*	*	*	*
EIGENA	23	21	20	21	22	20	20	20	20	20	20	20	20	20	20	20	20	20	20
EIGENA <sup>2</sup>	75	73	72	74	73	70	77	77	77	77	72	72	72	72	73	73	71	71	68
EIGENA <sup>3</sup>	92	92	92	91	91	86	92	92	92	92	92	92	92	92	92	93	92	91	93
EIGENB	171	140	143	150	135	200	147	147	147	147	179	172	156	149	141	151	152	151	161
EIGENB <sup>2</sup>	679	770	721	1085	681	1043	709	709	709	709	773	765	760	739	1233	1074	1006	960	788
EIGENC	65	68	65	63	66	64	67	67	67	67	65	65	65	65	64	74	68	66	64
$EIGENC^2$	245	275	256	253	234	287	275	275	275	275	260	260	259	257	274	298	245	230	232
GOTTFR	8	8	8	9	8	10	8	8	8	8	8	8	8	8	8	8	8	8	8
GROWTH	58	56	53	57	59	56	53	53	53	53	53	53	53	53	53	56	55	55	57
HATFLDF	28	27	26	25	25	39	26	26	26	26	25	25	25	25	35	25	25	25	25
HATFLDG	11	9	7	9	8	8	7	7	7	7	7	7	7	7	7	7	7	7	7
HEART6	164	164	159	160	163	275	159	159	163	163	159	159	159	159	164	162	162	162	162
HEART8	18	18	18	18	18	27	18	18	18	18	18	18	18	18	18	18	18	18	18

Table 5.3: Results for ACO-LS: parametric study for the interpolation rule (Algorithm 3.1).

ARC-LS

N. I. M. Gould, M. Porcelli and Ph. L. Toint

										ARC-L	5								
Name	0.001	0.005	$\beta$ 0.01	0.5	0.1	1	3.5	$\alpha_{max}$ 5	10	50	$10^{-12}$	$10^{-11}$	${\epsilon_{\chi} \over 10^{-10}}$	$10^{-9}$	$10^{-6}$	0.01	$\delta$ 0.05	$^{1}$ 0.25	0.5
HIMMELBA	4	4	4	5	5	5	4	4	4	4	4	4	4	4	4	4	4	4	4
HIMMELBC	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5
HIMMELBD	39	39	39	39	39	39	39	39	39	39	39	39	39	39	39	39	39	39	39
HIMMELBE	6	6	6	5	5	6	5	5	5	5	6	6	6	6	6	5	5	6	6
HS8	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6
HYDCAR6	460	461	437	490	470	493	448	448	448	448	434	472	466	449	455	399	408	448	462
HYPCIR	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5
INTEGREQ	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5
$INTEGREQ^2$	5	5	5	5	5	6	5	5	5	5	5	5	5	5	5	5	5	5	5
METHANB8	149	148	78	142	101	104	78	78	78	78	163	152	120	112	159	120	120	120	120
METHANL8	235	218	229	213	205	287	260	260	207	207	213	170	189	189	303	189	189	189	189
MSQRTA	40	40	40	40	40	35	40	35	35	35	40	40	38	40	44	39	39	39	37
$MSQRTA^2$	50	50	46	50	50	47	46	46	46	46	45	45	45	47	52	48	51	49	46
MSQRTB	38	38	38	39	39	35	38	34	34	34	38	38	40	38	42	38	38	39	35
$MSQRTB^2$	46	46	46	46	46	43	46	43	43	43	44	44	44	43	50	44	45	45	44
NYSTROM5	175	15	13	14	15	13	12	12	12	12	13	13	13	13	22	12	13	13	14
OSCIPANE	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7
PFIT1	160	155	150	176	168	152	2510	2510	2510	2510	150	150	150	150	1212	164	168	154	137
PFIT2	207	206	217	223	223	221	217	198	198	198	216	216	217	217	475	225	224	3349	221
PFIT3	272	295	280	287	272	276	280	4001	4001	4001	279	279	279	279	280	3880	284	318	297
PFIT4	379	381	406	367	395	367	373	373	411	386	403	403	403	404	412	412	385	380	383
POROUS1	36	35	32	34	37	42	32	32	32	32	32	32	32	32	32	32	32	32	32
POROUS1 <sup>2</sup>	78	76	80	83	93	82	80	80	80	80	80	80	80	80	80	80	80	81	81
POROUS1 <sup>3</sup>	175	164	175	169	95 168	161	186	186	186	186	175	175	173	175	175	181	175	173	175
POROUS1 POROUS2	48		74	109		>	74	78	146	146	74	175 74	74	74	74		1901	70	580
POROUS2 <sup>2</sup>		>			>											>			
	113	122	107	111	110	108	107	106	106	106	107	107	107	107	107	109	108	103	104
POROUS2 <sup>3</sup>	216	216	220	203	234	221	198	198	202	202	220	220	220	220	220	200	213	220	211
POROUS2 <sup>4</sup>	322	335	350	329	314	318	346	373	373	372	350	350	350	350	349	335	358	346	351
POWELLBS	169	192	181	180	160	168	181	181	181	181	90	88	99	133	528	99	99	99	99
POWELLSQ	13	13	13	13	13	13	417	417	417	417	13	13	13	13	13	13	13	13	13
QR3D	154	165	168	158	206	221	171	171	171	171	168	168	168	168	178	172	166	171	167
$QR3D^2$	1027	1019	1005	1022	1013	1041	1041	1041	1041	1041	956	954	954	969	1197	957	949	953	951
QR3DBD	64	71	65	63	70	70	65	65	65	65	64	64	64	64	65	64	64	64	64
$QR3DBD^2$	739	645	604	700	621	761	717	717	717	717	608	608	603	602	619	699	725	599	734
RECIPE	10	10	10	10	10	12	10	10	10	10	10	10	10	10	10	10	10	10	10
SINVALNE	23	23	23	23	23	21	23	23	23	23	23	23	23	23	23	23	23	23	23
SPMSQRT	12	12	12	12	12	12	12	12	12	14	12	12	12	12	12	12	12	13	14
TRIGGER	9	9	9	10	10	10	9	9	9	9	9	9	9	9	14	9	9	9	9
WOODSNE	28	28	28	28	28	29	28	28	28	28	28	28	28	28	28	28	28	28	28
YATP1SQ	47	47	43	48	41	46	43	43	43	43	43	43	43	43	43	43	43	43	43
$YATP1SQ^2$	23	23	23	23	23	24	23	23	23	23	23	23	23	23	23	23	23	23	23
YATP1SQ <sup>3</sup>	23	22	22	20	20	20	22	22	22	22	22	22	22	22	22	22	22	22	22
YATP2SQ	10	10	10	10	10	10	10	10	10	>	10	10	10	10	10	>	11	12	11
$YATP2SQ^2$	10	10	10	10	10	9	8	8	8	8	10	10	10	10	10	20	10	15	12
$YATP2SQ^3$	10	10	10	10	10	10	10	10	10	>	10	10	10	10	10	>	11	12	11
YFITNE	44	44	44	44	44	42	41	41	41	41	44	44	44	44	44	42	43	46	52
ZANGWIL3	5	5	5	6	6	7	5	5	4	4	5	5	5	5	5	5	5	6	7
	0	0	0	0	0	•	0	0	1	1	0	0	0	0	0	0	0	0	

Table 5.4: Results for ACO-LS: parametric study for the interpolation rule (Algorithm 3.1).

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ARC-LS

[]		$\delta_2$			r		I	ARC-LS	8	3			S	nax		g-rule
Name	0.25	$0.5^{-0.2}$	0.75	0.9	$(\eta_2 - \eta_1)/2$		1.25	1.5	2.5	3 3	4	10	50	500	1000	g-ruie
AIRCRFTA	4	4	4	4	$(\eta_2 - \eta_1)/2$	$\frac{\eta_2}{4}$	4	4	4	4	4	4	4	4	4	4
ARGAUSS	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3
ARGLALE	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6	15
ARGLALE	2	2	2	2	2	2	2	2	2	2	2	2	2	2	$\frac{0}{2}$	2
ARGTRIG	2 9	2 9	2	2 9	2 9	2 9	2 9	9	2 9	2 9	2 9	2 9	2 9	2 9	2 9	9
ARTIF	9 20	9 21	9 22	9 21	9 20	9 19	9 19	9 22	9 18	9 22	9 23	29	9 21	9 21	9 21	30
$ARTIF^2$	20 20	$\frac{21}{20}$	22	$\frac{21}{20}$	20 20	20	19 20	22	20	22	$\frac{23}{20}$	$\frac{29}{20}$	$\frac{21}{20}$	$\frac{21}{20}$	$\frac{21}{20}$	42
ARWDHNE					-			-			-	$20 \\ 258$				42 795
	SS	SS 40	SS 40	SS	ss 34	$\frac{ss}{34}$	$\frac{ss}{34}$	SS 24	SS 24	SS 24	$\frac{ss}{34}$	$\frac{258}{34}$	$\frac{ss}{34}$	$\frac{ss}{34}$	SS 24	
BDVALUES BDVALUES <sup>2</sup>	38 *	40 *	40 *	39 *	54 *	54 *	34 *	34 *	34 *	34 *	54 *	54 *	54 *	54 *	34 *	461 *
BOOTH																7
	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	
BRATU2D	6	6	6	6	6	6	6 6	6	6	6	6	6	6	6	6	12
BRATU2D <sup>2</sup>	6	6	6	6	6	6		6	6	6	6	6	6	6	6	24
BRATU2DT	14	14	14	14	14	14	14	15	14	14	14	14	14	14	14	29
BRATU2DT <sup>2</sup>	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	54
BRATU3D	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	13
BRATU3D <sup>2</sup>	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	24
BROWNALE	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4
BROWNALE <sup>2</sup> BROYDN3D	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3
	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8
BROYDN3D <sup>2</sup>	9	9	9	9	9	9	9	9	9	9	9	9	9	9	9	10
BROYDNBD	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
BROYDNBD <sup>2</sup>	10	11	11	11	11	11	11	11	11	11	11	11	11	11	11	12
CBRATU2D	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	20
CBRATU3D	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	16
CHANDHEQ	19	19	19	19	19	19	19	19	19	19	19	19	19	19	19	26
CHANNEL 2	159	141	133	129	142	151	164	146	141	148	146	146	145	139	139	237
CHANNEL <sup>2</sup> CHNRSBNE	111	98 42	91	89	105	102	107	99	98 40	99	102	101	96	94	97	214
	$\frac{46}{9}$	43 9	39 10	39 10	40	60 10	48	41	40	41	43	41	41	41	41	43 9
CLUSTER COOLHANS			10	10	10	10	10	10	10	10	10	10	10	10	10	~
CUBENE	416	600 14	$493 \\ 14$	684	441 14	441	441 14	441 14	$     441 \\     14 $	441	441 14	$491 \\ 14$	$435 \\ 13$	$934 \\ 14$	$934 \\ 14$	3403
DECONVNE	14	14		14		14	$14 \\ 15$	14 19		14	$14 \\ 16$					13
DRCAVTY1	$15 \\ 42$	$\frac{22}{43}$	16     42	19     47	19     45	$18 \\ 42$	15 49	19 43	$\frac{19}{43}$	$\frac{19}{43}$	10 42	19 61	$\frac{19}{47}$	$\frac{19}{40}$	19 40	30 53
DRCAVTY2	42	45 *	42	41 *	40	42	49 *	45	45 *	45 *	42	*	41 *	40 *	40 *	00 *
EIGENA	22	21	20	19	20	20	20	20	21	20	20	20	20	20	20	32
EIGENA <sup>2</sup>	$\frac{22}{75}$	65	20 62	19 71	20 72	$\frac{20}{75}$	$\frac{20}{77}$	20 71	21 74	$\frac{20}{73}$	$\frac{20}{74}$	$\frac{20}{75}$	$\frac{20}{72}$	$\frac{20}{72}$	$\frac{20}{72}$	260
EIGENA <sup>3</sup>	75 96	65 87	02 80	86	93	75 92	95	91	74 92	73 92	74 93	92	92	92	12 92	425
EIGENB	90 192	87 174	173		95 149	$\frac{92}{224}$	$\frac{95}{213}$	91 151	$\frac{92}{160}$	$\frac{92}{164}$	$93 \\ 172$	$\frac{92}{173}$	$\frac{92}{166}$	$\frac{92}{153}$	$\frac{92}{150}$	425 173
EIGENB <sup>2</sup>				155				$151 \\ 757$	$160 \\ 792$		172 822	$173 \\ 758$				1/3
EIGENC	1034	1318	998 64	806	868	818 107	2309 76			762 72			954	878 65	1149 62	
	73 225	78 200	64 204	68 272	73 258	107 272	76 568	67 276	95 272	73 260	65 272	68 248	80 250	65 220	62 224	67 462
EIGENC <sup>2</sup>	325 12	309	294	273	258	272 12	568	276	273	269	272 10	248	250	230	224	463
GOTTFR GROWTH	12 52	9 55	8 55	9 56	9 52	12	12 186	8	8 54	9 50	10 86	8 52	8 52	8 52	8 52	16 157
HATFLDF	$53 \\ 26$	$\frac{55}{22}$	$\frac{55}{24}$	$\frac{56}{25}$	$53 \\ 25$	$93 \\ 26$	186     27		$\frac{54}{25}$	$\frac{52}{25}$	$\frac{86}{25}$	$53 \\ 25$	$\frac{53}{25}$	$\frac{53}{25}$	$\frac{53}{25}$	157 19
HATFLDG	$\frac{20}{7}$	22 7	$\frac{24}{7}$	$\frac{25}{7}$	25 7	$\frac{20}{7}$	21 7	25 7	$\frac{25}{7}$	$\frac{25}{7}$	$\frac{25}{7}$	$\frac{25}{7}$	$\frac{25}{7}$	$\frac{25}{7}$	$\frac{25}{7}$	19
HEART6	165	144	$\frac{7}{151}$	158	158	171	171	164	164	168	7 169	132	7 159	7 159	159	554
HEART8	165 28	144 28	$\frac{151}{23}$	158 24	158	171	171	164 18	164 18	168 18	169 18	132 18	159 18	159 18	159 18	554 21
IILAN10	20	20	20	24	10	10	10	10	10	10	10	10	10	10	10	21

Table 5.5: Results for ACO-LS: parametric study for the interpolation rule (Algorithm 3.1) and g-rule (3.5) (last column).

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		δ	2		1	1	-	no-Lo	δ	3			$\delta_m$	ar		g-rule
Name	0.25	0.5	0.75	0.9	$(\eta_2 - \eta_1)/2$	$\eta_2$	1.25	1.5	2.5	3	4	10	50	500	1000	0
HIMMELBA	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	8
HIMMELBC	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5
HIMMELBD	38	38	40	40	39	29	23	57	33	29	26	40	39	40	40	51
HIMMELBE	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6	8
HS8	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6
HYDCAR6	569	450	443	512	424	488	484	461	460	447	431	494	480	485	486	534
HYPCIR	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5
INTEGREQ	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5
INTEGREQ <sup>2</sup>	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	6
METHANB8	81	156	124	129	120	120	120	120	120	120	120	120	120	120	120	81
METHANL8	263	241	152	224	188	181	203	191	186	194	283	307	202	212	200	391
MSQRTA	38	38	38	38	40	40	40	40	40	40	40	40	40	40	40	53
MSQRTA <sup>2</sup>	53	49	46	44	51	50	50	45	45	45	45	45	45	45	45	64
MSQRTB	36	36	37	36	38	38	38	38	38	38	38	38	38	38	38	53
MSQRTB <sup>2</sup>	44	40	41	42	45	46	46	44	44	44	44	45	43	44	44	59
NYSTROM5	11	40 12	12	12	13	13	13	13	13	13	13	13	13	13	13	14
OSCIPANE	7	7	7	12	7	7	15 7	7	13	7	7	7	7	7	7	7
PFIT1	159	152	155	166	141	165	975	150	165	150	176	141	167	152	164	148
PFIT2	224	$132 \\ 212$	213	212	213	213	975 213	212	210	228	260	220	223	$\frac{152}{235}$	230	148
PFIT3	224 289	212 284	$213 \\ 260$	$212 \\ 272$	213 279	$213 \\ 271$	1892	264	$\frac{210}{291}$	327	$\frac{200}{317}$	3965	223 289	235 285	$\frac{230}{315}$	>
PFIT4	400	284 396	200 399	401	403	403	403	368	391	383	483	400	400	372	399	>
POROUS1	400 34	390 39	399 35	39	403 37	$\frac{403}{34}$	$\frac{403}{37}$	32	33	32	$\frac{463}{39}$	400 36	400	572	399 40	51
POROUS1 <sup>2</sup>	- 34 - 88	39 81	33 83	39 95	57 74	34 86	88	-32 78	33 83	32 82	39 77	- 30 - 84	$\frac{40}{78}$	$\frac{52}{70}$	$\frac{40}{75}$	111
POROUS1 <sup>3</sup>	00 167	179	85 171		195		152	190	$^{00}$	82 181	192	84 195	78 179	163	162	*
POROUS1 POROUS2	226	179	>	171 >	>	164 >	$152 \\ 210$	>	>	>	192 90	>	>	105	-	61
POROUS2 POROUS2 <sup>2</sup>	220 97	114	125	110	128	127	116	114	104	112	115	222	111	119	$\frac{ss}{129}$	125
POROUS2 <sup>3</sup>	230	219	$125 \\ 227$	188	223	$\frac{127}{216}$	226	218	$104 \\ 229$	$112 \\ 197$	224	222	204	212	$129 \\ 240$	$123 \\ 271$
POROUS2 <sup>4</sup>	230 333	328	343	322	223 341	210 *	352	323	303	367	224 334	220 *	$\frac{204}{349}$	335	3240	467
POWELLBS	- 335 - 96	328 90	545 86	322 87	99	99	352 99	323 99	303 99	307 99	334 99	99	549 99	335 99	524 99	407 1314
POWELLSQ	90 13	90 13	80 13		99 12	99 13		99 17		99 17		99 13	99 13	99 13		1514 16
•				13			14		18		16				13	
QR3D	339	236	157	165	173	180	164	160	153	161	161	168	168	168	168	205
QR3D <sup>2</sup>	1538	1344	1069	1016	942	969 62	859	919 60	946 70	968	957	950 64	954	954	954	947
QR3DBD QR3DBD <sup>2</sup>	91	80	68 71.2	63 607	64 760	63 722	64	69 69	70 670	63 797	63 710	64 667	64	64	64	122
RECIPE	868	954	713	697	769	733	640 10	686	670 10	737	716	667	604	593	593	728
	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	15
SINVALNE	23	22	21	21	23	25 19	37	23	24	23	23	21	23	23	23	26
SPMSQRT	12	12	12	12	12	12	12	12	12	12	12	12	12	12	12	29
TRIGGER	9	9	9	9	9	9	9	9	9	9	9	9	9	9	9	10
WOODSNE VATED1SO	23	26	27	28	28	28	28	28	28	28	28 46	28	28	28	28	41
YATP1SQ	41	46	43	43	46	45	52	43	43	43	46	49	47	41	43	39
YATP1SQ <sup>2</sup>	23	23	23	23	24	25 22	22	23	23	23	23	26	23	23	23	21
YATP1SQ <sup>3</sup>	22	22	22	22	23	22	19 10	22	22	22	22	23	22	22	22	20
YATP2SQ	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	14
YATP2SQ <sup>2</sup>	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	14
YATP2SQ <sup>3</sup>	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	14
YFITNE	44	44	44	44	44	44	44	44	44	44	44	45	40	41	43	453
ZANGWIL3	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	23

Table 5.6: Results for ACO-LS: parametric study for the interpolation rule (Algorithm 3.1) and g-rule (3.5) (last column).