

A Multi-preconditioned GMRES Algorithm

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Joint work with:

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The specific problem we are interested in:

Consider solving

$$Ax = b$$

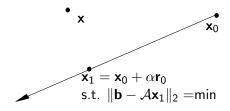
where ${\cal A}$ is a large, sparse matrix, using a Krylov subspace method.

Suppose we have two (or more!) possible preconditioners,

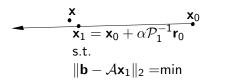
$$\mathcal{P}_1$$
 and \mathcal{P}_2 ,

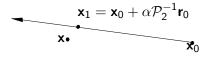
Can we (optimally) combine information from more than one preconditioner?













Multi-preconditioning



Two relevant methods

Multi-preconditioned conjugate gradients (MPCG): Bridson & Greif (2006)

- Combines multiple preconditioners automatically in a (locally) optimal way
- ▶ Requires A and $\{P_i\}$ to be symmetric positive definite
- We lose the short-term-recurrence of PCG

Flexible GMRES (FGMRES): Saad (1993)

- ▶ Allows variable preconditioners e.g. \mathcal{P}_1 on odd iterations, \mathcal{P}_2 on even iterations
- Uses all preconditioners, but nontrivial subspace is being constructed and optimality properties are not fully understood



Complete Multi-preconditioned Arnoldi

```
Pick \mathbf{x}_0, let V_1 = \mathbf{r}_0 / ||\mathbf{r}_0||.
Let Z_1 = [\mathcal{P}_1^{-1}V_1 \cdots \mathcal{P}_t^{-1}V_1] \in \mathbb{R}^{n \times t}
for i = 1 \dots \max its
       Q = AZ
       for i = 1 \dots i
              H_{i,i} = V_i^T W
               W = W - V_i H_{i,i}
       end
       W = V_{i+1}H_{i+1,i} (skinny QR factorization)
       Z_{i+1} = [\mathcal{P}_1^{-1} V_{i+1} \cdots \mathcal{P}_t^{-1} V_{i+1}]
end
```

Complete Multi-preconditioned Arnoldi

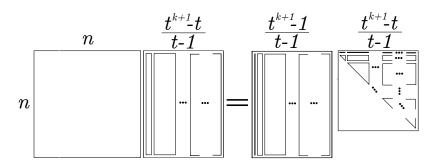
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```

$$\mathcal{A}[Z_1 \cdots Z_k] = [V_1 \dots V_{k+1}]\widetilde{H}_k$$



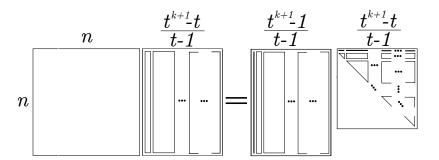
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Complete Multipreconditioned Arnoldi

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Impractical!

Selective Multi-preconditioned Arnoldi

```
Pick \mathbf{x}_0, let V_1 = \mathbf{r}_0 / \|\mathbf{r}_0\|.
Let Z_1 = [\mathcal{P}_1^{-1}V_1 \cdots \mathcal{P}_t^{-1}V_1] \in \mathbb{R}^{n \times t}
for i = 1...
       Q = AZ_i
       for i = 1 \dots i
               H_{j,i} = V_j^T W
               W = W - V_i H_{ii}
       end
        W = V_{i+1}H_{i+1,i} (skinny QR factorization)
       Z_{i+1} = [\mathcal{P}_1^{-1} V_{i+1}^{(1)} \cdots \mathcal{P}_t^{-1} V_{i+1}^{(t)}]
end
```

Selective Multi-preconditioned Arnoldi

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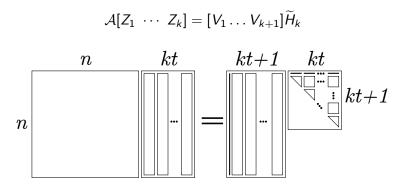
Various possibilities for generating the search directions

Could replace $Z_{i+1} = [\mathcal{P}_1^{-1}V_{i+1}^{(1)}\cdots\mathcal{P}_t^{-1}V_{i+1}^{(t)}]$ by taking a mix of all columns, say: $Z_{i+1} = [\mathcal{P}_1^{-1}V_{i+1}\mathbf{1}\cdots\mathcal{P}_t^{-1}V_{i+1}\mathbf{1}],$ where $\mathbf{1}$ is a vector of all ones.

Practical evidence shows "mixing" is typically more effective; no analytical observations to support this.



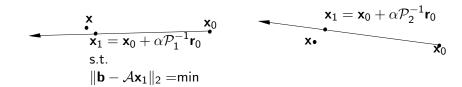
Selective Multi-preconditioned Arnoldi



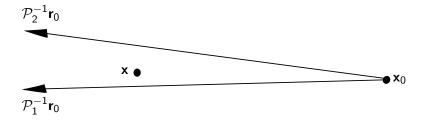
Multi-preconditioned GMRES

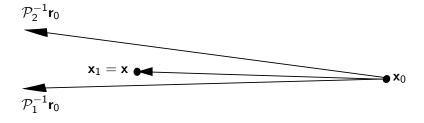
Find vector \mathbf{y}_k s.t. $\mathbf{x}_k = \mathbf{x}_0 + Z_k \mathbf{y}_k$ i.e., find \mathbf{y}_k which minimizes

$$\begin{aligned} \|\mathbf{b} - \mathcal{A}\mathbf{x}_{k}\|_{2} &= \|\mathbf{b} - \mathcal{A}(\mathbf{x}_{0} + [Z_{1} \cdots Z_{k}]\mathbf{y}_{k})\|_{2} \\ &= \|\mathbf{r}_{0} - \mathcal{A}[Z_{1} \cdots Z_{k}]\mathbf{y}_{k}\|_{2} \\ &= \|\mathbf{r}_{0} - [V_{1} \cdots V_{k+1}]\widetilde{H}_{k}\mathbf{y}_{k}\|_{2} \\ &= \|V_{1}\|\mathbf{r}_{0}\|_{2} - [V_{1} \cdots V_{k+1}]\widetilde{H}_{k}\mathbf{y}_{k}\|_{2} \\ &= \|[V_{1} \cdots V_{k+1}](\|\mathbf{r}_{0}\|_{2}\mathbf{e}_{1} - \widetilde{H}_{k}\mathbf{y}_{k})\|_{2} \\ &= \|\|\mathbf{r}_{0}\|_{2}\mathbf{e}_{1} - \widetilde{H}_{k}\mathbf{y}_{k}\|_{2} \end{aligned}$$









Comparison of costs

Costs at the kth iteration:

	Matvec. prods	inner products	pre. solves
MPGMRES	t ^k	$\frac{t^{2k+1}+t^{2k}+t^{k+1}-3t^k}{2(t-1)}$	t ^k
tMPGMRES	t	$(k-\frac{1}{2})t^{2'}+\frac{3}{2}t$	t
GMRES	1	$\bar{k}{+}1$	1



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GMRES	1	k+1	1

Easily parallelized



Characterizing the search space

GMRES preconditioned with a right-preconditioner \mathcal{P} finds the vector that minimizes the 2-norm of the residual over all vectors of the form

$$\mathbf{x}^{(k)} = \mathbf{x}^{(0)} + \mathcal{P}^{-1}\mathbf{y}_k,$$

where \mathbf{y}_k is a member of the Krylov subspace

$$\mathcal{K}_k(\mathcal{AP}^{-1}, \mathbf{r}^{(0)}) = \operatorname{span}(\mathbf{r}^{(0)}, \mathcal{AP}^{-1}\mathbf{r}^{(0)}, \dots, (\mathcal{AP}^{-1})^{k-1}\mathbf{r}^{(0)}).$$



Characterizing the search space

The extension for (complete) MPGMRES: the first two iterates satisfy

$$\textbf{x}^{(1)} - \textbf{x}^{(0)} \in \mathrm{span}\{\mathcal{P}_1^{-1}\textbf{r}^{(0)}, \mathcal{P}_2^{-1}\textbf{r}^{(0)}\}$$

$$\begin{split} \textbf{x}^{(2)} - \textbf{x}^{(0)} &\in \operatorname{span}\{\mathcal{P}_{1}^{-1}\textbf{r}^{(0)}, \mathcal{P}_{2}^{-1}\textbf{r}^{(0)}, \mathcal{P}_{1}^{-1}\mathcal{A}\mathcal{P}_{1}^{-1}\textbf{r}^{(0)}, \mathcal{P}_{1}^{-1}\mathcal{A}\mathcal{P}_{2}^{-1}\textbf{r}^{(0)}, \\ \mathcal{P}_{2}^{-1}\mathcal{A}\mathcal{P}_{1}^{-1}\textbf{r}^{(0)}, \mathcal{P}_{2}^{-1}\mathcal{A}\mathcal{P}_{2}^{-1}\textbf{r}^{(0)}\}, \end{split}$$

and the rest follow the same pattern.



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$$\begin{split} \mathbf{x}^{(1)} - \mathbf{x}^{(0)} &\in \operatorname{span}\{\mathcal{P}_{1}^{-1}\mathbf{r}^{(0)}, \mathcal{P}_{2}^{-1}\mathbf{r}^{(0)}\} \\ &= \operatorname{span}\{\boldsymbol{X}\mathbf{e}^{(0)}, \boldsymbol{Y}\mathbf{e}^{(0)}\} \\ \mathbf{x}^{(2)} - \mathbf{x}^{(0)} &\in \operatorname{span}\{\mathcal{P}_{1}^{-1}\mathbf{r}^{(0)}, \mathcal{P}_{2}^{-1}\mathbf{r}^{(0)}, \mathcal{P}_{1}^{-1}\mathcal{A}\mathcal{P}_{1}^{-1}\mathbf{r}^{(0)}, \mathcal{P}_{1}^{-1}\mathcal{A}\mathcal{P}_{2}^{-1}\mathbf{r}^{(0)}, \\ &\qquad \qquad \mathcal{P}_{2}^{-1}\mathcal{A}\mathcal{P}_{1}^{-1}\mathbf{r}^{(0)}, \mathcal{P}_{2}^{-1}\mathcal{A}\mathcal{P}_{2}^{-1}\mathbf{r}^{(0)}\}, \\ &= \operatorname{span}\{\boldsymbol{X}\mathbf{e}^{(0)}, \boldsymbol{Y}\mathbf{e}^{(0)}, \boldsymbol{X}^{2}\mathbf{e}^{(0)}, \boldsymbol{X}\boldsymbol{Y}\mathbf{e}^{(0)}\boldsymbol{Y}\boldsymbol{X}\mathbf{e}^{(0)}, \boldsymbol{Y}^{2}\mathbf{e}^{(0)}\} \\ &\qquad \qquad \boldsymbol{X} = \mathcal{P}_{1}^{-1}\mathcal{A}, \ \boldsymbol{Y} = \mathcal{P}_{2}^{-1}\mathcal{A} \end{split}$$

and the rest follow the same pattern.

Demonstration of the relative richness of the search space

Given \mathcal{P}_1 and \mathcal{P}_2 , take $\mathbf{x}^{(0)} = \mathbf{0}$. Then

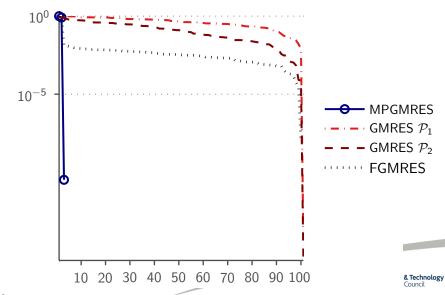
$$\mathcal{P}_1^{-1}\mathcal{A}\mathcal{P}_2^{-1}\mathbf{b}$$

lies in the search space after two iterations.

Therefore if **b** is an eigenvector of $\mathcal{AP}_1^{-1}\mathcal{AP}_2^{-1}$, MPGMRES will converge after two iterations.



Two iterations



Breakdowns

All breakdowns in standard GMRES are 'lucky'.

This is not the case with MPGMRES....

e.g. if $\mathcal{P}_1 = \mathcal{P}_2$, the matrix $Z_1 = [\mathcal{P}_1^{-1} \mathbf{r}^{(0)} \ \mathcal{P}_2^{-1} \mathbf{r}^{(0)}]$ will be of rank one

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...but not a problem in general – we can monitor the subdiagonal entries of the upper Hessenberg matrix:

- if 0 on subdiagonal, and not converged must be a linearly dependent vector: discard and all is fine
- ▶ if 0 on the subdiagonal, and residual small enough lucky breakdown!
- ▶ if no zero on subdiagonal, no problem!



Examples



Domain decomposition

Consider the advection-diffusion equation on $\Omega = [0,1]^2$:

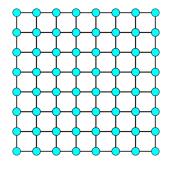
$$-\nabla^2 u + \omega \cdot \nabla u = f \quad \text{in } \Omega$$
$$u = 0 \quad \text{on } \partial \Omega.$$

Upon discretization by finite differences we get the matrix equation

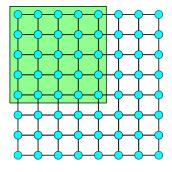
$$A\mathbf{u} = \mathbf{b},$$

where A is a real positive, but nonsymmetric, matrix.

Domain decomposition (cont.)

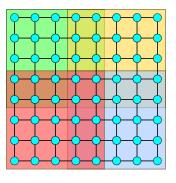


Domain decomposition (cont.)





Domain decomposition (cont.)



$$A^{(i)} = R_{i,\delta} A R_{i,\delta}^T, i = 1, ..., 4,$$

where $R_{i,\delta}$ is a restriction matrix, and δ denotes the number of nodes overlapping ($\delta = 1$ here).



Additive Schwarz preconditioner

The additive Schwarz preconditioner has its inverse defined as

$$M^{-1} = \sum_{i=1}^{t} R_{i,\delta}^{T} (A^{(i)})^{-1} R_{i,\delta}.$$



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The additive Schwarz preconditioner has its inverse defined as

$$M^{-1} = \sum_{i=1}^t R_{i,\delta}^T (A^{(i)})^{-1} R_{i,\delta}.$$

Well suited to a multi-preconditioned approach: take each solve on a subdomain as a preconditioner.



Restricted Additive Schwarz

The restricted Additive Schwarz preconditioner is defined as

$$M^{-1} = \sum_{i=1}^{t} R_{i,0}^{T} (R_{i,\delta} A R_{i,\delta}^{T})^{-1} R_{i,\delta}.$$

This has the effect of removing the overlap in the preconditioner, hence improving convergence.

Also an ideal candidate for multi-preconditioning.

Two subdomains

In the special case where we have two subdomains without overlap, we can show that complete and selective MPGMRES are equivalent.



Two subdomains

In the special case where we have two subdomains without overlap, we can show that complete and selective MPGMRES are equivalent.

Ν	sMPGMRES	GMRES	Ratio
2^{2}	$6.59 \times 10^{-3} (5)$	$9.79 \times 10^{-3} (9)$	0.67
2^3	1.86×10^{-2} (8)	1.36×10^{-2} (12)	1.37
2^{4}	2.79×10^{-2} (11)	4.14×10^{-2} (17)	0.67
2^{5}	1.08×10^{-1} (16)	1.50×10^{-1} (24)	0.72
2^{6}	4.66×10^{-1} (19)	7.37×10^{-1} (33)	0.63
2^{7}	2.78 (25)	4.85 (46)	0.57
2^{8}	17.1 (30)	34.2 (65)	0.50

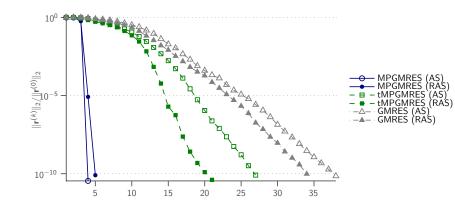


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A PDE constrained optimization problem

$$\min_{y,u} \frac{1}{2} ||y - \hat{y}||_2^2 + \frac{\beta}{2} ||u||_2^2$$

s.t.
$$-\nabla^2 y = u \text{ in } \Omega$$

 $y = f \text{ on } \partial \Omega$



Two preconditioners

$$\begin{bmatrix} \beta Q & 0 & -Q \\ 0 & Q & K \\ -Q & K & 0 \end{bmatrix} \begin{bmatrix} \mathbf{u} \\ \mathbf{y} \\ \mathbf{p} \end{bmatrix} = \begin{bmatrix} 0 \\ \mathbf{b} \\ \mathbf{d} \end{bmatrix}$$
a choice of preconditioner

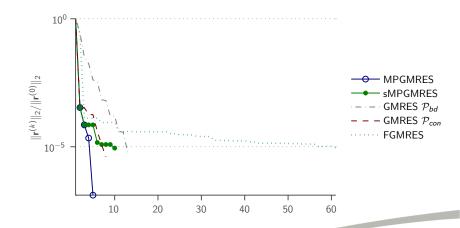
$$\mathcal{P}_1 = \left[egin{array}{cccc} eta Q & 0 & 0 \ 0 & Q & 0 \ 0 & 0 & KQ^{-1}K \end{array}
ight]$$
block diagonal

$$\mathcal{P}_1 = \begin{bmatrix} \beta Q & 0 & 0 \\ 0 & Q & 0 \\ 0 & 0 & KQ^{-1}K \end{bmatrix} \qquad \mathcal{P}_2 = \begin{bmatrix} 0 & 0 & -Q \\ 0 & \beta KQ^{-1}K & K \\ -Q & K & 0 \end{bmatrix}$$
block diagonal constraint

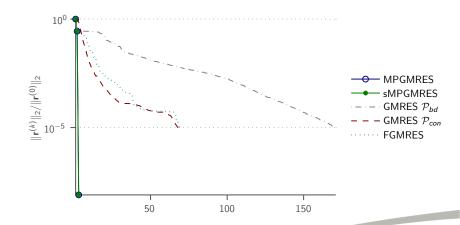
[Rees, Dollar & Wathen (2010)]



 $\beta = 10^{-4}$



$$\beta = 10^{-8}$$



Timings

	$\beta = 10^{-4}$	$\beta = 10^{-8}$
complete MPGMRES	10.8	2.1
selective MPGMRES	12.3	1.4
GMRES, \mathcal{P}_{bd}	3.5	53.5
GMRES, \mathcal{P}_{con}	3.1	28.4
FGMRES	26.4	33.7



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MPGMRES is an extension of the standard preconditioned GMRES which allows us to use more than one preconditioner.

The method:

- seems to work well when we have non-ideal preconditioners which complement each other
- can handle any number of candidate preconditioners
- can be parallelized, obtaining potential computational gains

Paper and MATLAB code available at

www.numerical.rl.ac.uk/people/rees/

Fortran 95 code (HSL_MI28) under development