# A Multi-preconditioned GMRES Algorithm 

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## The specific problem we are interested in:

Consider solving

$$
\mathcal{A} \mathbf{x}=\mathbf{b}
$$

where $\mathcal{A}$ is a large, sparse matrix, using a Krylov subspace method.
Suppose we have two (or more!) possible preconditioners,

$$
\mathcal{P}_{1} \quad \text { and } \quad \mathcal{P}_{2},
$$

Can we (optimally) combine information from more than one preconditioner?

## Diagrams



## Diagrams


s.t.

$$
\left\|\mathbf{b}-\mathcal{A} \mathbf{x}_{1}\right\|_{2}=\min
$$



# Multi-preconditioning 

## Two relevant methods

Multi-preconditioned conjugate gradients (MPCG):
Bridson \& Greif (2006)

- Combines multiple preconditioners automatically in a (locally) optimal way
- Requires $\mathcal{A}$ and $\left\{\mathcal{P}_{i}\right\}$ to be symmetric positive definite
- We lose the short-term-recurrence of PCG

Flexible GMRES (FGMRES): Saad (1993)

- Allows variable preconditioners - e.g. $\mathcal{P}_{1}$ on odd iterations, $\mathcal{P}_{2}$ on even iterations
- Uses all preconditioners, but nontrivial subspace is being constructed and optimality properties are not fully understood


## Complete Multi-preconditioned Arnoldi

```
Pick \(\mathbf{x}_{0}\), let \(V_{1}=\mathbf{r}_{0} /\left\|\mathbf{r}_{0}\right\|\).
Let \(Z_{1}=\left[\mathcal{P}_{1}^{-1} V_{1} \cdots \mathcal{P}_{t}^{-1} V_{1}\right] \in \mathbb{R}^{n \times t}\)
for \(i=1 \ldots\) max_its
    \(Q=A Z_{i}\)
    for \(j=1 \ldots i\)
    \(H_{j, i}=V_{j}^{T} W\)
    \(W=W-V_{j} H_{j, i}\)
    end
    \(W=V_{i+1} H_{i+1, i} \quad\) (skinny QR factorization)
    \(Z_{i+1}=\left[\begin{array}{llll}\mathcal{P}_{1}^{-1} & V_{i+1} & \cdots & \mathcal{P}_{t}^{-1} V_{i+1}\end{array}\right]\)
end
```


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Let $Z_{1}=\left[\mathcal{P}_{1}^{-1} V_{1} \cdots \mathcal{P}_{t}^{-1} V_{1}\right] \in \mathbb{R}^{n \times t}$
for $i=1 \ldots$ max_its

$$
\begin{aligned}
& Q=A Z_{i} \\
& \text { for } j=1 \ldots i \\
& \quad H_{j, i}=V_{j}^{\top} W \\
& \text { end } W=W-V_{j} H_{j, i}
\end{aligned}
$$

$$
W=V_{i+1} H_{i+1, i} \quad \text { (skinny QR factorization) }
$$

$$
Z_{i+1}=\left[\begin{array}{llll}
\mathcal{P}_{1}^{-1} V_{i+1} & \cdots & \mathcal{P}_{t}^{-1} V_{i+1}
\end{array}\right]
$$

end

$$
\mathcal{A}\left[Z_{1} \ldots Z_{k}\right]=\left[V_{1} \ldots V_{k+1}\right] \widetilde{H}_{k}
$$

## Complete Multipreconditioned Arnoldi

$$
\mathcal{A}\left[z_{1} \cdots Z_{k}\right]=\left[\begin{array}{lll}
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$$



Impractical!

## Selective Multi-preconditioned Arnoldi

$$
\begin{aligned}
& \text { Pick } \mathbf{x}_{0} \text {, let } V_{1}=\mathbf{r}_{0} /\left\|\mathbf{r}_{0}\right\| \text {. } \\
& \text { Let } Z_{1}=\left[\mathcal{P}_{1}^{-1} V_{1} \cdots \mathcal{P}_{t}^{-1} V_{1}\right] \in \mathbb{R}^{n \times t} \\
& \text { for } i=1 \ldots \\
& \quad Q=A Z_{i} \\
& \quad \text { for } j=1 \ldots i \\
& \quad H_{j, i}=V_{j}^{T} W \\
& \quad W=W-V_{j} H_{j, i} \\
& \quad \text { end } \\
& \quad W=V_{i+1} H_{i+1, i} \quad(\text { skinny QR factorization) } \\
& \quad Z_{i+1}=\left[\mathcal{P}_{1}^{-1} V_{i+1}^{(1)} \cdots \mathcal{P}_{t}^{-1} V_{i+1}^{(t)}\right]
\end{aligned}
$$

## Selective Multi-preconditioned Arnoldi

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Let $Z_{1}=\left[\mathcal{P}_{1}^{-1} V_{1} \cdots \mathcal{P}_{t}^{-1} V_{1}\right] \in \mathbb{R}^{n \times t}$
for $i=1 \ldots$
$Q=A Z_{i}$
for $j=1 \ldots i$

$$
\begin{aligned}
& H_{j, i}=V_{j}^{\top} W \\
& W=W-V_{j} H_{j, i}
\end{aligned}
$$

end

$$
W=V_{i+1} H_{i+1, i} \quad \text { (skinny QR factorization) }
$$

$$
Z_{i+1}=\left[\begin{array}{llll}
\mathcal{P}_{1}^{-1} V_{i+1}^{(1)} & \cdots & \mathcal{P}_{t}^{-1} V_{i+1}^{(t)}
\end{array}\right]
$$

end

$$
\mathcal{A}\left[Z_{1} \ldots Z_{k}\right]=\left[\begin{array}{lll}
V_{1} & \ldots & V_{k+1}
\end{array}\right] \widetilde{H}_{k}
$$

## Various possibilities for generating the search directions

Could replace $Z_{i+1}=\left[\mathcal{P}_{1}^{-1} V_{i+1}^{(1)} \cdots \mathcal{P}_{t}^{-1} V_{i+1}^{(t)}\right]$ by taking a mix of all columns, say: $Z_{i+1}=\left[\mathcal{P}_{1}^{-1} V_{i+1} 1 \cdots \mathcal{P}_{t}^{-1} V_{i+1} 1\right]$, where $\mathbf{1}$ is a vector of all ones.

Practical evidence shows "mixing" is typically more effective; no analytical observations to support this.

## Selective Multi-preconditioned Arnoldi

$$
\mathcal{A}\left[z_{1} \cdots z_{k}\right]=\left[v_{1} \ldots v_{k+1}\right] \tilde{H}_{k}
$$



## Multi-preconditioned GMRES

Find vector $\mathbf{y}_{k}$ s.t. $\mathbf{x}_{k}=\mathbf{x}_{0}+Z_{k} \mathbf{y}_{k}$
i.e., find $y_{k}$ which minimizes

$$
\begin{aligned}
\left\|\mathbf{b}-\mathcal{A} \mathbf{x}_{k}\right\|_{2} & =\left\|\mathbf{b}-\mathcal{A}\left(\mathbf{x}_{0}+\left[Z_{1} \cdots Z_{k}\right] \mathbf{y}_{k}\right)\right\|_{2} \\
& =\left\|\mathbf{r}_{0}-\mathcal{A}\left[Z_{1} \cdots Z_{k}\right] \mathbf{y}_{k}\right\|_{2} \\
& =\left\|\mathbf{r}_{0}-\left[V_{1} \cdots V_{k+1}\right] \widetilde{H_{k}} \mathbf{y}_{k}\right\|_{2} \\
& =\left\|V_{1}\right\| \mathbf{r}_{0}\left\|_{2}-\left[V_{1} \cdots V_{k+1}\right] \widetilde{H}_{k} \mathbf{y}_{k}\right\|_{2} \\
& =\left\|\left[V_{1} \cdots V_{k+1}\right]\left(\left\|\mathbf{r}_{0}\right\|_{2} \mathbf{e}_{1}-\widetilde{H_{k}} \mathbf{y}_{k}\right)\right\|_{2} \\
& =\| \| \mathbf{r}_{0}\left\|_{2} \mathbf{e}_{1}-\widetilde{H_{k}} \mathbf{y}_{k}\right\|_{2}
\end{aligned}
$$

## Diagrams


s.t.

$$
\left\|\mathbf{b}-\mathcal{A} \mathbf{x}_{1}\right\|_{2}=\min
$$

## Diagrams



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## Comparison of costs

Costs at the kth iteration:

|  | Mat.-vec. prods | inner products | pre. solves |
| :---: | :---: | :---: | :---: |
| MPGMRES | $t^{k}$ | $\frac{t^{2 k+1}+t^{2 k}+t^{k+1}-3 t^{k}}{2(t-1)}$ | $t^{k}$ |
| tMPGMRES | $t$ | $\left(k-\frac{1}{2}\right) t^{2}+\frac{3}{2} t$ | $t$ |
| GMRES | 1 | $\mathrm{k}+1$ | 1 |

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## Characterizing the search space

GMRES preconditioned with a right-preconditioner $\mathcal{P}$ finds the vector that minimizes the 2 -norm of the residual over all vectors of the form

$$
\mathbf{x}^{(k)}=\mathbf{x}^{(0)}+\mathcal{P}^{-1} \mathbf{y}_{k}
$$

where $\mathbf{y}_{k}$ is a member of the Krylov subspace

$$
\mathcal{K}_{k}\left(\mathcal{A} \mathcal{P}^{-1}, \mathbf{r}^{(0)}\right)=\operatorname{span}\left(\mathbf{r}^{(0)}, \mathcal{A} \mathcal{P}^{-1} \mathbf{r}^{(0)}, \ldots,\left(\mathcal{A} \mathcal{P}^{-1}\right)^{k-1} \mathbf{r}^{(0)}\right)
$$

## Characterizing the search space

The extension for (complete) MPGMRES: the first two iterates satisfy

$$
\begin{aligned}
& \mathbf{x}^{(1)}-\mathbf{x}^{(0)} \in \operatorname{span}\left\{\mathcal{P}_{1}^{-1} \mathbf{r}^{(0)}, \mathcal{P}_{2}^{-1} \mathbf{r}^{(0)}\right\} \\
& \mathbf{x}^{(2)}-\mathbf{x}^{(0)} \in \operatorname{span}\left\{\mathcal{P}_{1}^{-1} \mathbf{r}^{(0)}, \mathcal{P}_{2}^{-1} \mathbf{r}^{(0)}, \mathcal{P}_{1}^{-1} \mathcal{A} \mathcal{P}_{1}^{-1} \mathbf{r}^{(0)}, \mathcal{P}_{1}^{-1} \mathcal{A} \mathcal{P}_{2}^{-1} \mathbf{r}^{(0)},\right. \\
&\left.\mathcal{P}_{2}^{-1} \mathcal{A} \mathcal{P}_{1}^{-1} \mathbf{r}^{(0)}, \mathcal{P}_{2}^{-1} \mathcal{A} \mathcal{P}_{2}^{-1} \mathbf{r}^{(0)}\right\},
\end{aligned}
$$

and the rest follow the same pattern.

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&= \operatorname{span}\left\{X \mathbf{e}^{(0)}, Y \mathbf{e}^{(0)}\right\} \\
& \mathbf{x}^{(2)}-\mathbf{x}^{(0)} \in \operatorname{span}\left\{\mathcal{P}_{1}^{-1} \mathbf{r}^{(0)}, \mathcal{P}_{2}^{-1} \mathbf{r}^{(0)}, \mathcal{P}_{1}^{-1} \mathcal{A} \mathcal{P}_{1}^{-1} \mathbf{r}^{(0)}, \mathcal{P}_{1}^{-1} \mathcal{A} \mathcal{P}_{2}^{-1} \mathbf{r}^{(0)},\right. \\
&\left.\mathcal{P}_{2}^{-1} \mathcal{A} \mathcal{P}_{1}^{-1} \mathbf{r}^{(0)}, \mathcal{P}_{2}^{-1} \mathcal{A} \mathcal{P}_{2}^{-1} \mathbf{r}^{(0)}\right\}, \\
&= \operatorname{span}\left\{X \mathbf{e}^{(0)}, Y \mathbf{e}^{(0)}, X^{2} \mathbf{e}^{(0)}, X Y \mathbf{e}^{(0)} Y X \mathbf{e}^{(0)}, Y^{2} \mathbf{e}^{(0)}\right\} \\
& X=\mathcal{P}_{1}^{-1} \mathcal{A}, Y=\mathcal{P}_{2}^{-1} \mathcal{A}
\end{aligned}
$$

and the rest follow the same pattern.

## Demonstration of the relative richness of the search space

Given $\mathcal{P}_{1}$ and $\mathcal{P}_{2}$, take $\mathbf{x}^{(0)}=\mathbf{0}$. Then

$$
\mathcal{P}_{1}^{-1} \mathcal{A} \mathcal{P}_{2}^{-1} \mathbf{b}
$$

lies in the search space after two iterations.
Therefore if $\mathbf{b}$ is an eigenvector of $\mathcal{A} \mathcal{P}_{1}^{-1} \mathcal{A} \mathcal{P}_{2}^{-1}$, MPGMRES will converge after two iterations.

## Two iterations



$$
\begin{aligned}
& \text { - - MPGMRES } \\
& \text { - - - GMRES } \mathcal{P}_{1} \\
& \text { - - - GMRES } \mathcal{P}_{2} \\
& \text { FGMRES }
\end{aligned}
$$

## Breakdowns

All breakdowns in standard GMRES are 'lucky'.
This is not the case with MPGMRES....
e.g. if $\mathcal{P}_{1}=\mathcal{P}_{2}$, the matrix $Z_{1}=\left[\mathcal{P}_{1}^{-1} \mathbf{r}^{(0)} \mathcal{P}_{2}^{-1} \mathbf{r}^{(0)}\right]$ will be of rank one

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...but not a problem in general - we can monitor the subdiagonal entries of the upper Hessenberg matrix:

- if 0 on subdiagonal, and not converged - must be a linearly dependent vector: discard and all is fine
- if 0 on the subdiagonal, and residual small enough - lucky breakdown!
- if no zero on subdiagonal, no problem!


## Examples

## Domain decomposition

Consider the advection-diffusion equation on $\Omega=[0,1]^{2}$ :

$$
\begin{aligned}
-\nabla^{2} u+\omega \cdot \nabla u & =f & & \text { in } \Omega \\
u & =0 & & \text { on } \partial \Omega .
\end{aligned}
$$

Upon discretization by finite differences we get the matrix equation

$$
A \mathbf{u}=\mathbf{b}
$$

where $A$ is a real positive, but nonsymmetric, matrix.

## Domain decomposition (cont.)



## Domain decomposition (cont.)



## Domain decomposition (cont.)



$$
A^{(i)}=R_{i, \delta} A R_{i, \delta}^{T}, i=1, \ldots, 4
$$

where $R_{i, \delta}$ is a restriction matrix, and $\delta$ denotes the number of nodes overlapping ( $\delta=1$ here).

## Additive Schwarz preconditioner

The additive Schwarz preconditioner has its inverse defined as

$$
M^{-1}=\sum_{i=1}^{t} R_{i, \delta}^{T}\left(A^{(i)}\right)^{-1} R_{i, \delta}
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$$

Well suited to a multi-preconditioned approach: take each solve on a subdomain as a preconditioner.

## Restricted Additive Schwarz

The restricted Additive Schwarz preconditioner is defined as

$$
M^{-1}=\sum_{i=1}^{t} R_{i, 0}^{T}\left(R_{i, \delta} A R_{i, \delta}^{T}\right)^{-1} R_{i, \delta}
$$

This has the effect of removing the overlap in the preconditioner, hence improving convergence.

Also an ideal candidate for multi-preconditioning.

## Two subdomains

In the special case where we have two subdomains without overlap, we can show that complete and selective MPGMRES are equivalent.

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| N | sMPGMRES | GMRES | Ratio |
| :---: | :---: | :---: | :---: |
| $2^{2}$ | $6.59 \times 10^{-3}(5)$ | $9.79 \times 10^{-3}(9)$ | 0.67 |
| $2^{3}$ | $1.86 \times 10^{-2}(8)$ | $1.36 \times 10^{-2}(12)$ | 1.37 |
| $2^{4}$ | $2.79 \times 10^{-2}(11)$ | $4.14 \times 10^{-2}(17)$ | 0.67 |
| $2^{5}$ | $1.08 \times 10^{-1}(16)$ | $1.50 \times 10^{-1}(24)$ | 0.72 |
| $2^{6}$ | $4.66 \times 10^{-1}(19)$ | $7.37 \times 10^{-1}(33)$ | 0.63 |
| $2^{7}$ | $2.78(25)$ | $4.85(46)$ | 0.57 |
| $2^{8}$ | $17.1(30)$ | $34.2(65)$ | 0.50 |

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## A PDE constrained optimization problem

$$
\begin{aligned}
& \qquad \begin{aligned}
\min _{y, u} \frac{1}{2}\|y-\hat{y}\|_{2}^{2} & +\frac{\beta}{2}\|u\|_{2}^{2} \\
\text { s.t. } \quad-\nabla^{2} y & =u \text { in } \Omega \\
y & =f \text { on } \partial \Omega
\end{aligned}
\end{aligned}
$$

## Two preconditioners

$$
\begin{aligned}
& {\left[\begin{array}{ccc}
\beta Q & 0 & -Q \\
0 & Q & K \\
-Q & K & 0
\end{array}\right]\left[\begin{array}{l}
\mathbf{u} \\
\mathbf{y} \\
\mathbf{p}
\end{array}\right]=\left[\begin{array}{l}
0 \\
\mathbf{b} \\
\mathbf{d}
\end{array}\right]} \\
& \text { a choice of preconditioner } \\
& \mathcal{P}_{1}=\left[\begin{array}{ccc}
\beta Q & 0 & 0 \\
0 & Q & 0 \\
0 & 0 & K Q^{-1} K
\end{array}\right] \\
& \text { block diagonal } \\
& \mathcal{P}_{2}=\left[\begin{array}{ccc}
0 & 0 & -Q \\
0 & \beta K Q^{-1} K & K \\
-Q & K & 0
\end{array}\right] \\
& \text { constraint }
\end{aligned}
$$

[Rees, Dollar \& Wathen (2010)]

## $\beta=10^{-4}$




## Timings

|  | $\beta=10^{-4}$ | $\beta=10^{-8}$ |
| :---: | :---: | :---: |
| complete MPGMRES | 10.8 | 2.1 |
| selective MPGMRES | 12.3 | 1.4 |
| GMRES, $\mathcal{P}_{b d}$ | 3.5 | 53.5 |
| GMRES, $\mathcal{P}_{c o n}$ | 3.1 | 28.4 |
| FGMRES | 26.4 | 33.7 |

## Timings

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MPGMRES is an extension of the standard preconditioned GMRES which allows us to use more than one preconditioner.
The method:

- seems to work well when we have non-ideal preconditioners which complement each other
- can handle any number of candidate preconditioners
- can be parallelized, obtaining potential computational gains

Paper and MATLAB code available at
www.numerical.rl.ac.uk/people/rees/

Fortran 95 code (HSL_MI28) under development

